

**COURSE:
THEORY OF
AUTOMATA
COMPUTATION**

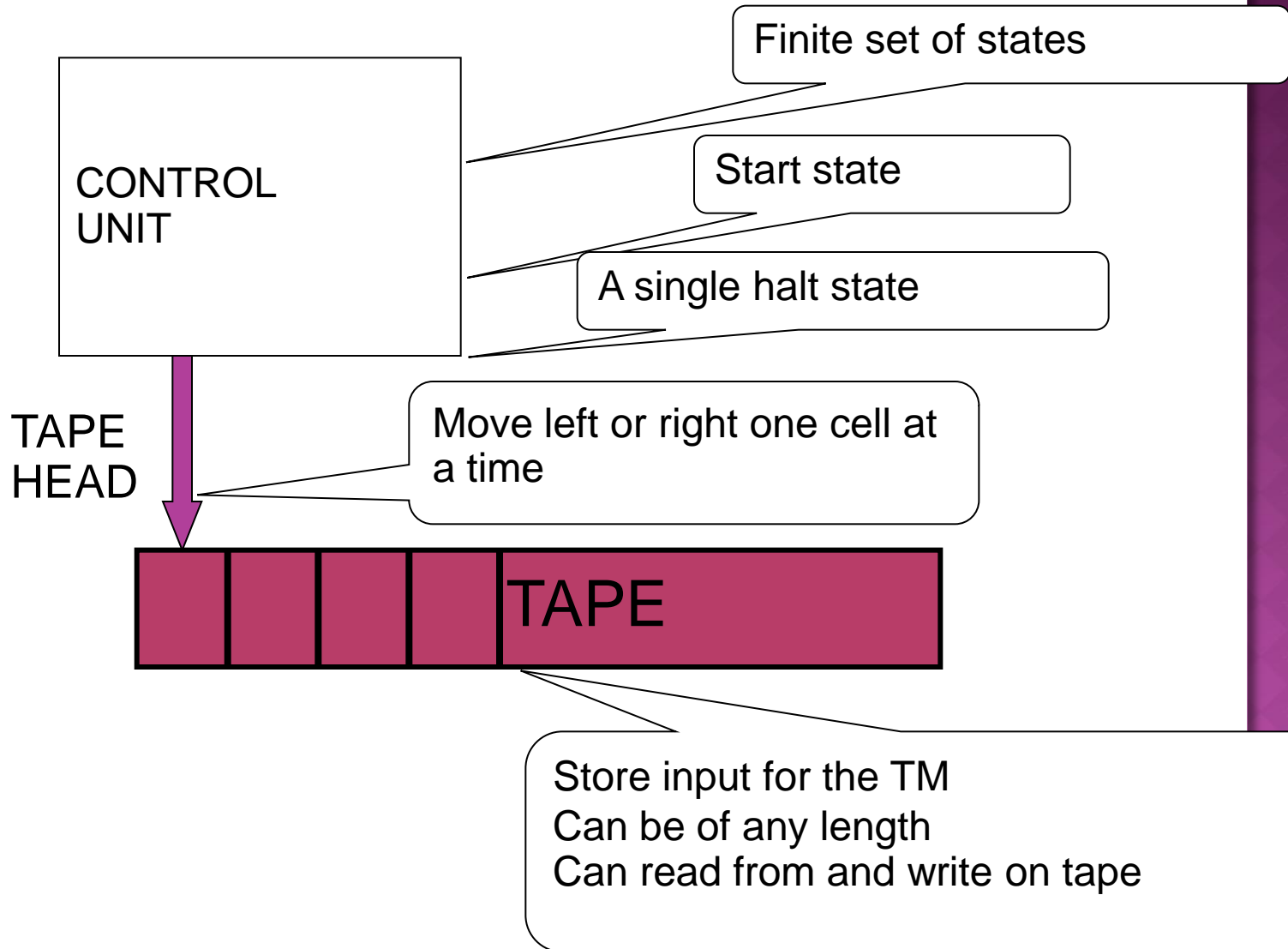
TOPICS TO BE COVERED

- ◉ Turing Machines (TM)
- ◉ Model of Computation

OUTLINES

- ◉ Structure of Turing machines
- ◉ Deterministic Turing machines (DTM)
 - Accepting a language
 - Computing a function
- ◉ Composite Turing machines
- ◉ Multitape Turing machines
- ◉ Nondeterministic Turing machines (NTM)
- ◉ Universal Turing machines (UTM)

STRUCTURE OF TM



WHAT DOES A TM DO?

- ⦿ Determine if an input x is in a language.
 - That is, answer if the answer of a problem P for the instance x is “yes”.
- ⦿ Compute a function
 - Given an input x , what is $f(x)$?

HOW DOES A TM WORK?

- ⦿ At the beginning,
 - A TM is in the *start state (initial state)*
 - its tape head points at the first cell
 - The tape contains Δ , following by input string, and the rest of the tape contains Δ .

HOW DOES A TM WORK?

- ◎ For each move, a TM
 - reads the symbol under its tape head
 - According to the *transition function* on the symbol read from the tape and its current state, the TM:
 - write a symbol on the tape
 - move its tape head to the left or right one cell or not
 - changes its state to the *next state*

WHEN DOES A TM STOP WORKING?

- ⦿ A TM stops working,
 - when it gets into the special state called **halt state**. (**halts**)
 - The output of the TM is on the tape.
 - when the tape head is on the leftmost cell and is moved to the left. (**hangs**)
 - when there is no **next state**. (**hangs**)

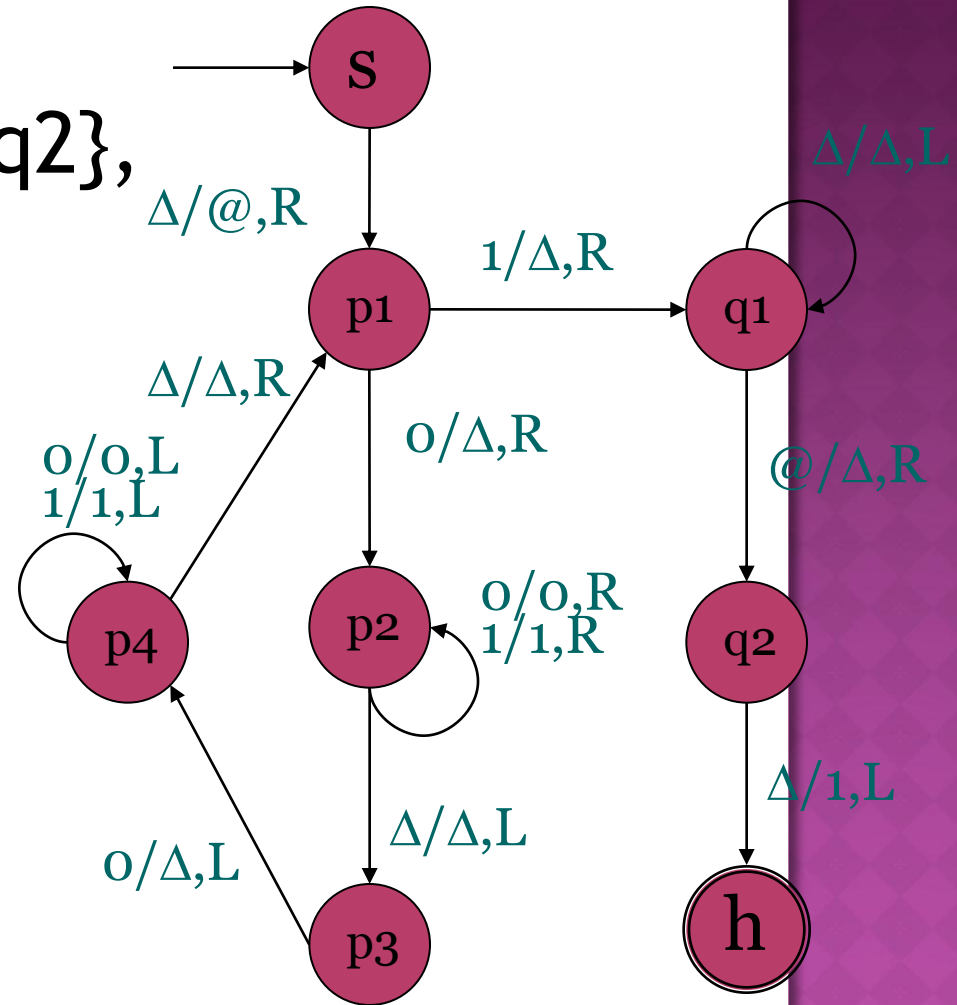
HOW TO DEFINE DETERMINISTIC TM (DTM)

- a quintuple $(Q, \Sigma, \Gamma, \delta, s)$, where
 - the set of states Q is finite, not containing halt state h ,
 - the input alphabet Σ is a finite set of symbols not including the blank symbol Δ ,
 - the tape alphabet Γ is a finite set of symbols containing Σ , but not including the blank symbol Δ ,
 - the start state s is in Q , and
 - the transition function δ is a partial function from $Q \times (\Gamma \cup \{\Delta\}) \rightarrow Q \cup \{h\} \times (\Gamma \cup \{\Delta\}) \times \{L, R, S\}$.

EXAMPLE OF A DTM

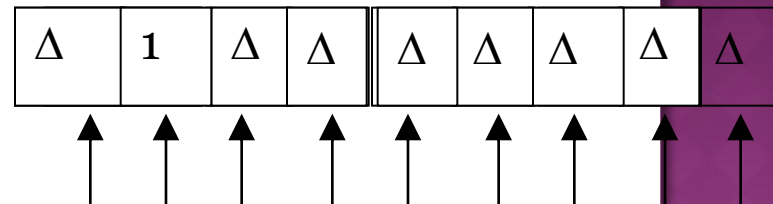
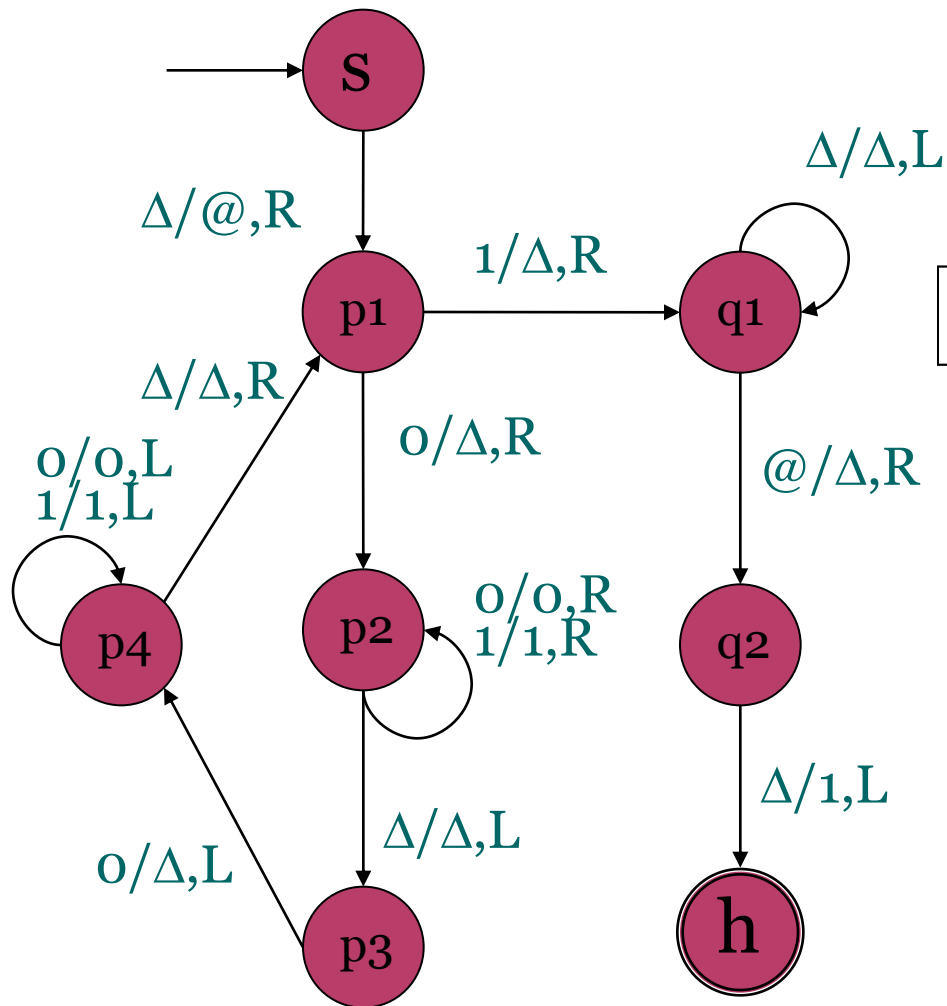
⊙ $M =$

$(\{s, p1, p2, p3, p4, q1, q2\},$
 $\{0, 1\}, \{0, 1, @, \delta, s\}$



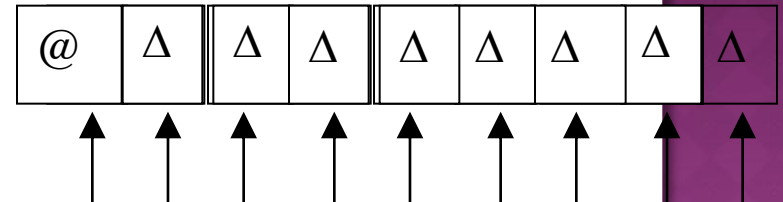
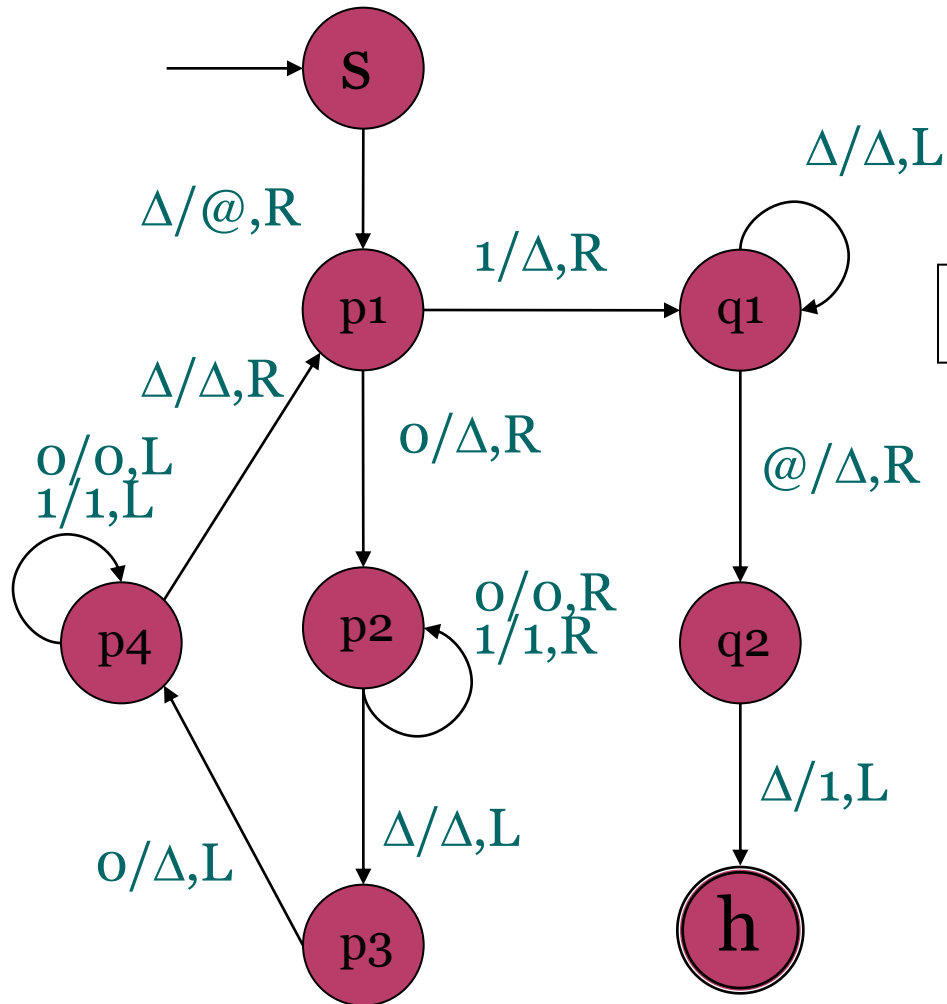


HOW A DTM WORKS



On the input 0001000, the TM halts.

HOW A DTM WORKS



On the input 0000000 , the TM hangs.

CONFIGURATION

Definition

Let $T = (Q, \Sigma, \Gamma, \delta, s)$ be a DTM.

A configuration of T is an element of

$$Q \cup \{h\} \times (\Gamma \cup \{\Delta\})^* \times (\Gamma \cup \{\Delta\})$$

$$\times (\Gamma \cup \{\Delta\})^*$$

Current state

as

string to the right of tape head

symbol under tape head

string to the left of tape head

■ $(q, l \cdot \underline{a} \cdot r)$

YIELD THE NEXT CONFIGURATION

Definition

- Let $T = (Q, \Sigma, \Gamma, \delta, s)$ be a DTM, and $(q_1, \alpha_1 \underline{a}_1 \beta_1)$ and $(q_2, \alpha_2 \underline{a}_2 \beta_2)$ be two configurations of T .

We say $(q_1, \alpha_1 \underline{a}_1 \beta_1)$ **yields** $(q_2, \alpha_2 \underline{a}_2 \beta_2)$ **in one step**, denoted by $(q_1, \alpha_1 \underline{a}_1 \beta_1) \vdash^T (q_2, \alpha_2 \underline{a}_2 \beta_2)$, if

- $\delta(q_1, a_1) = (q_2, a_2, s)$, $\alpha_1 = \alpha_2$ and $\beta_1 = \beta_2$,
- $\delta(q_1, a_1) = (q_2, b, R)$, $\alpha_2 = \alpha_1 b$ and $\beta_1 = a_2 \beta_2$,

YIELD IN ZERO STEP OR MORE

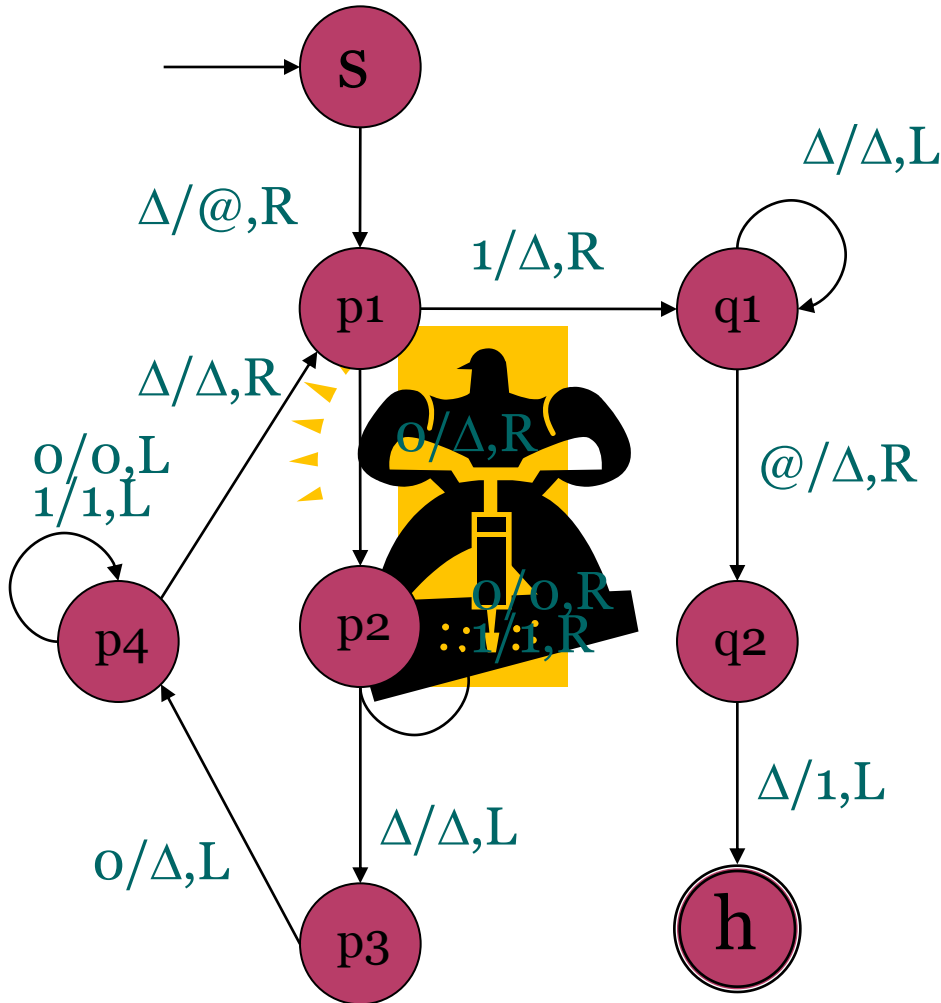
Definition

- Let $T=(Q, \Sigma, \Gamma, \delta, s)$ be a DTM, and $(q_1, \alpha_1 \underline{a}_1 \beta_1)$ and $(q_2, \alpha_2 \underline{a}_2 \beta_2)$ be two configurations of T .

We say $(q_1, \alpha_1 \underline{a}_1 \beta_1)$ **yields** $(q_2, \alpha_2 \underline{a}_2 \beta_2)$ **in zero step or more**, denoted by $(q_1, \alpha_1 \underline{a}_1 \beta_1) \vdash_T^* (q_2, \alpha_2 \underline{a}_2 \beta_2)$, if

- $q_1=q_2, \alpha_1=\alpha_2, \underline{a}_1=\underline{a}_2,$ and $\beta_1=\beta_2$, or
- $(q_1, \alpha_1 \underline{a}_1 \beta_1) \vdash_T (q, \alpha \underline{a} \beta)$ and $(q, \alpha \underline{a} \beta) \vdash_T^* (q_2, \alpha_2 \underline{a}_2 \beta_2)$ for some q in Q , α and β in Γ^* ,

YIELD IN ZERO STEP OR MORE: EXAMPLE



$(s, \underline{\Delta}0001000)$
 $(p1, @\underline{0}001000)$
 $(p2, @\underline{\Delta}001000)$
 $(p2, @\underline{\Delta}001000\underline{\Delta})$
 $(p3, @\underline{\Delta}001000\underline{\quad})$
 $(p4, @\underline{\Delta}00100\underline{\Delta})$
 $(p4, @\underline{\Delta}00100\underline{\Delta})$
 $(p1, @\underline{\Delta}00100\underline{\Delta})$
 $(p2, @\underline{\Delta}\underline{0}100\underline{\Delta})$
 $(p2, @\underline{\Delta}\underline{0}100\underline{\Delta})$
 $(p3, @\underline{\Delta}\underline{0}100\underline{\quad})$

$(p4, @\underline{\Delta}\underline{\Delta}010\underline{\quad})$
 $(p4, @\underline{\Delta}\underline{\Delta}010\underline{\quad})$
 $(p1, @\underline{\Delta}\underline{\Delta}010\underline{\quad})$
 $(p2, @\underline{\Delta}\underline{\Delta}\underline{\Delta}10\underline{\quad})$
 $(p2, @\underline{\Delta}\underline{\Delta}\underline{\Delta}10\underline{\quad})$
 $(p2, @\underline{\Delta}\underline{\Delta}\underline{\Delta}10\underline{\Delta})$
 $(p3, @\underline{\Delta}\underline{\Delta}\underline{\Delta}10\underline{\quad})$
 $(p4, @\underline{\Delta}\underline{\Delta}\underline{\Delta}1\underline{\quad})$
 $(p4, @\underline{\Delta}\underline{\Delta}\underline{\Delta}1\underline{\quad})$
 $(p1, @\underline{\Delta}\underline{\Delta}\underline{\Delta}1\underline{\quad})$
 $(p2, @\underline{\Delta}\underline{\Delta}\underline{0}100\underline{\Delta})$
 $(q1, @\underline{\Delta}\underline{\Delta}\underline{\Delta}\underline{\Delta}\underline{\Delta})$
 $(q1, @)$
 $(q2, \underline{\Delta}\underline{\Delta})$
 $(h, \underline{\Delta}1)$