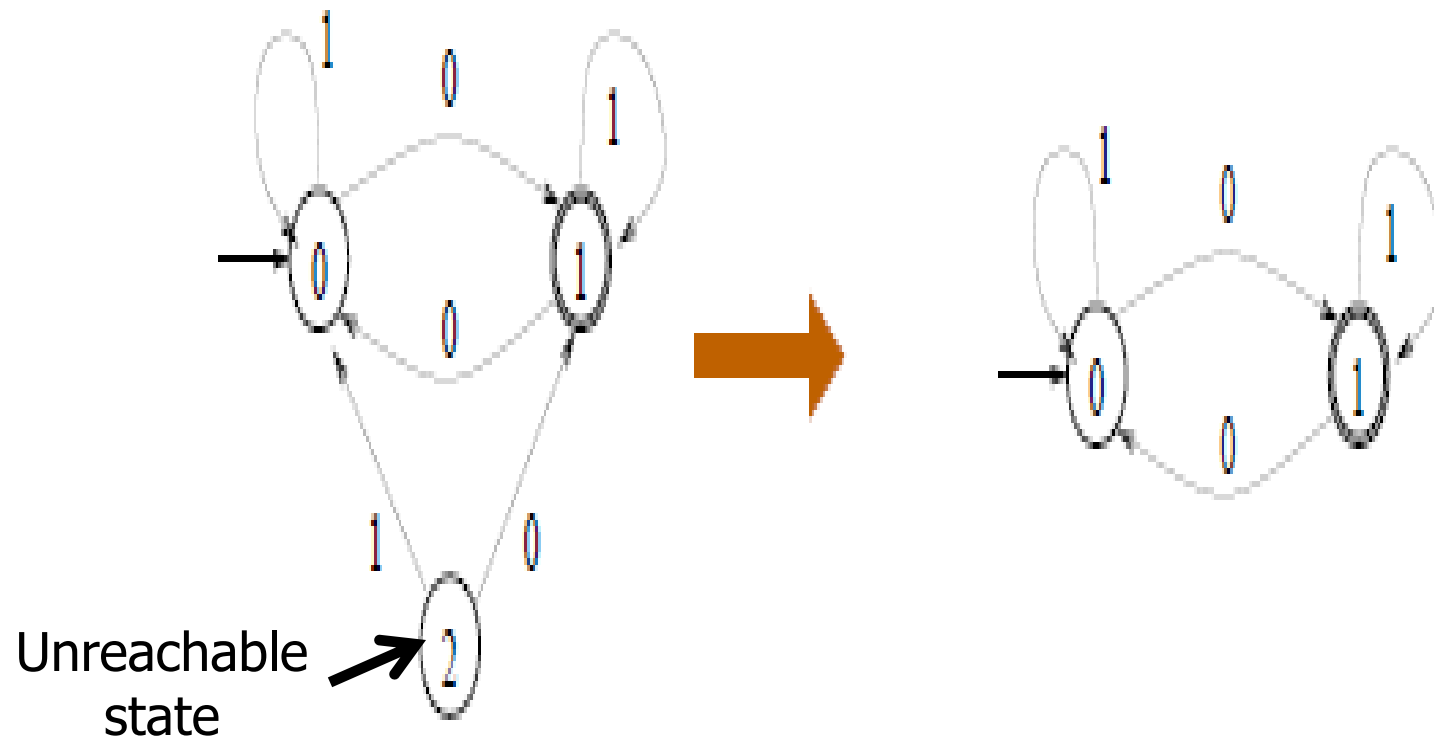


**COURSE:  
THEORY OF  
AUTOMATA  
COMPUTATION**

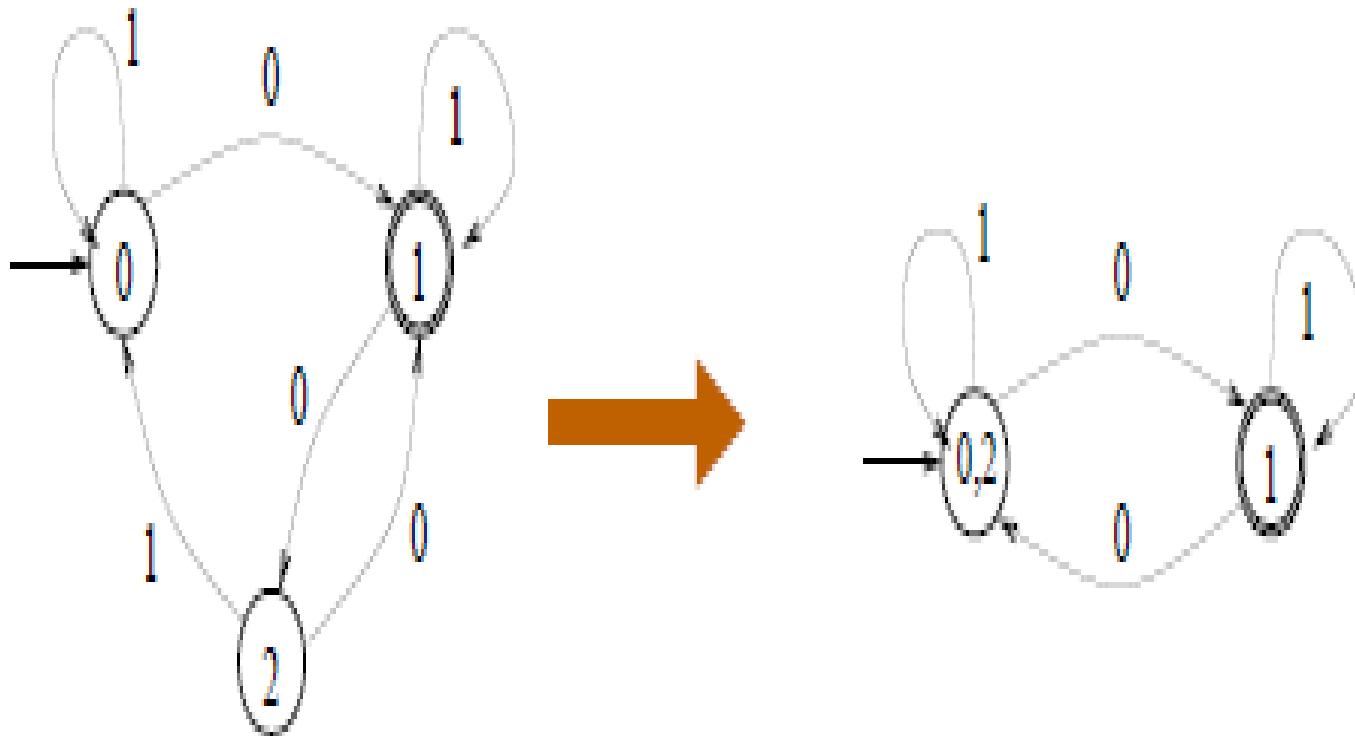
# TOPICS TO BE COVERED

- ◉ Minimization of FA
- ◉ Minimization Algorithm
  - Guarantees smallest possible DFA for a given regular language
  - Proof of this fact.

# EXAMPLE ONE



# EXAMPLE TWO

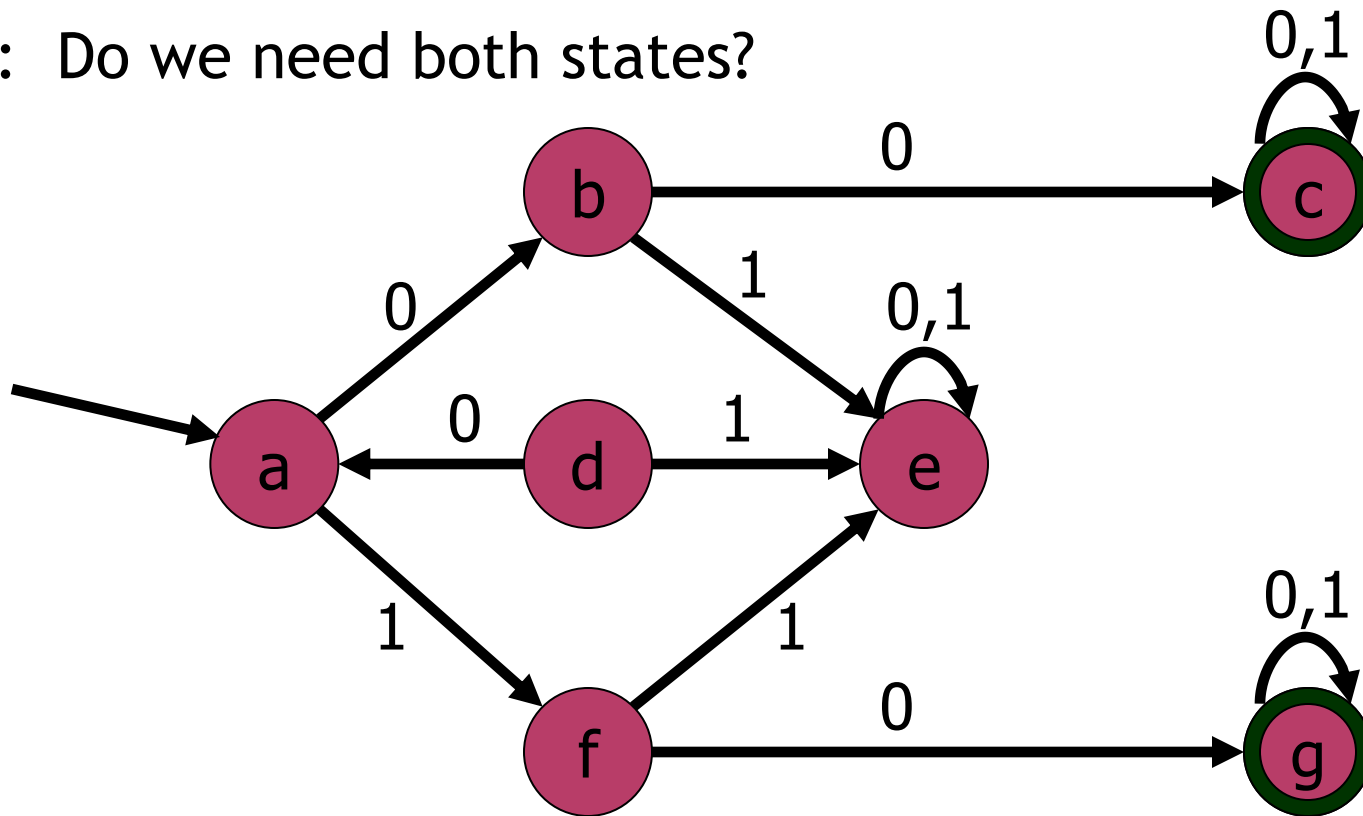


# EQUIVALENT STATES.

## EXAMPLE

Consider the accept states  $c$  and  $g$ . They are both sinks meaning that any string which ever reaches them is guaranteed to be accepted later.

Q: Do we need both states?

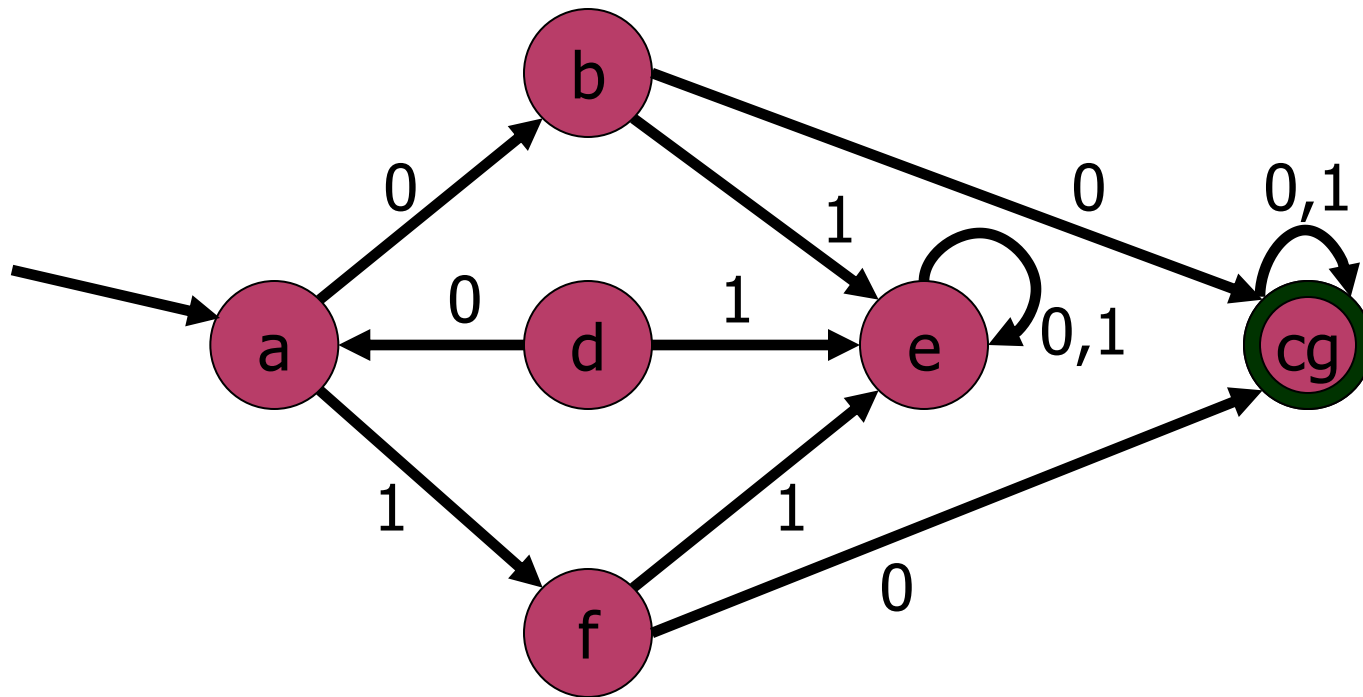


# EQUIVALENT STATES.

## EXAMPLE

A: No, they can be unified as illustrated below.

Q: Can any other states be unified because any subsequent string suffixes produce identical results?



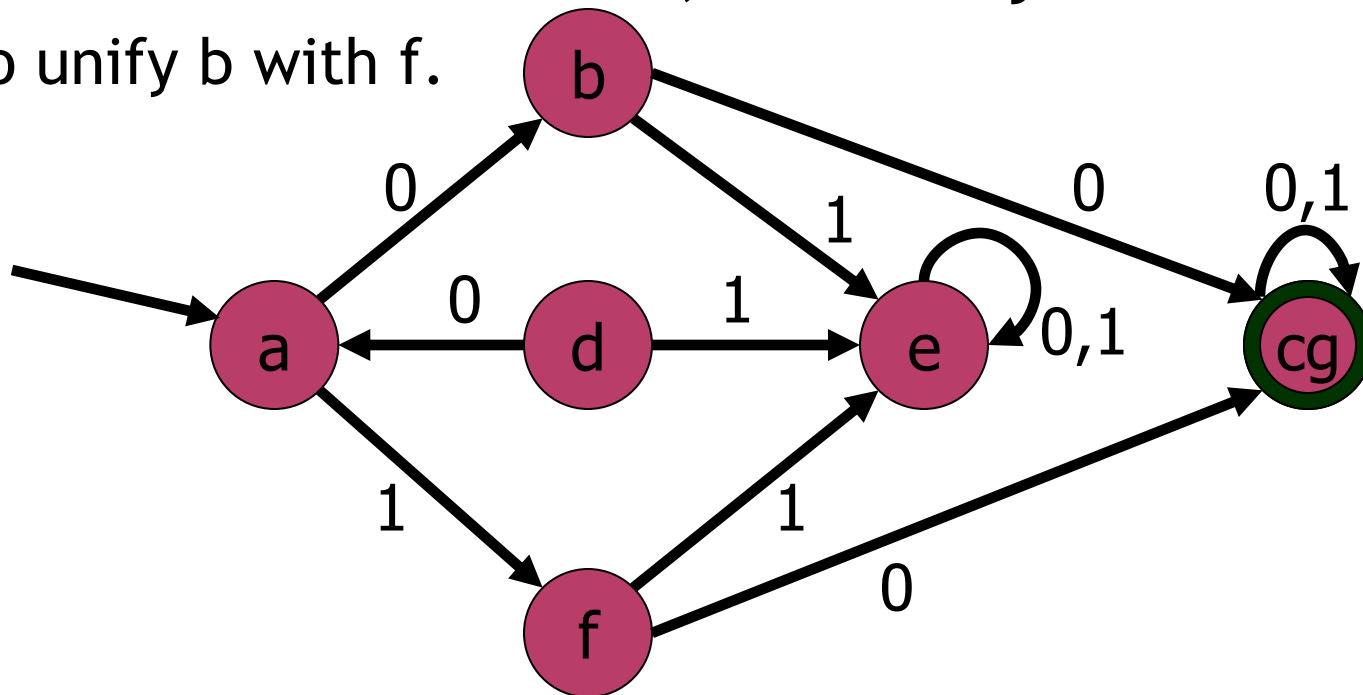
# EQUIVALENT STATES.

## EXAMPLE

A: Yes, b and f. Notice that if you're in b or f then:

1. if string ends, reject in both cases
2. if next character is 0, forever accept in both cases
3. if next character is 1, forever reject in both cases

So unify b with f.

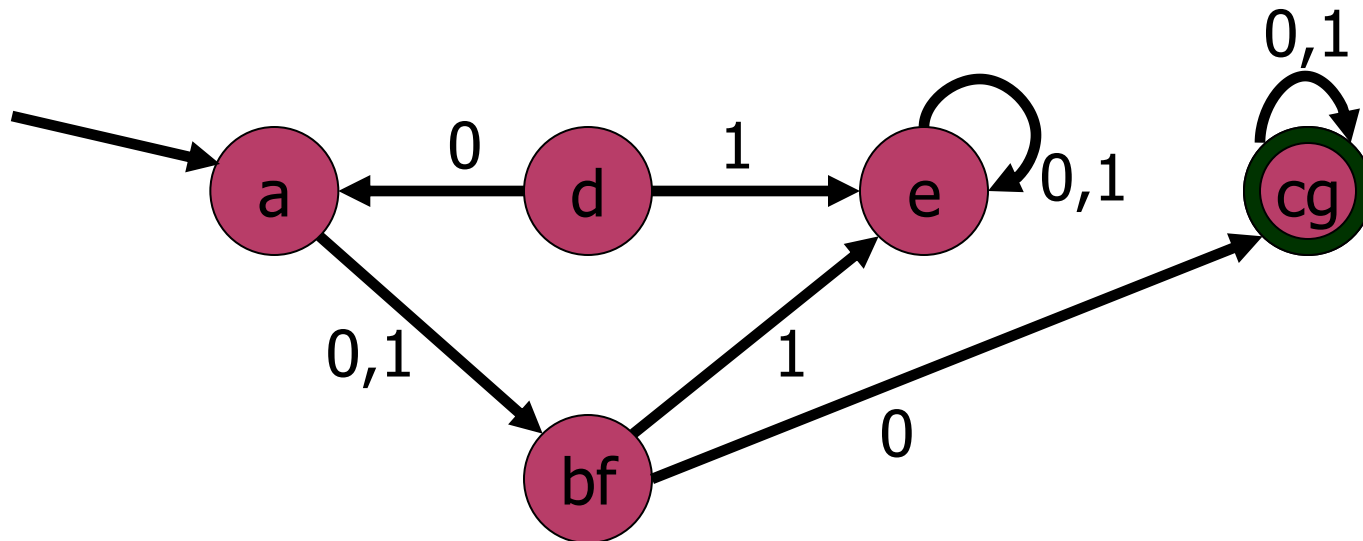


# EQUIVALENT STATES.

## EXAMPLE

Intuitively two states are equivalent if all subsequent behavior from those states is the same.

Q: Come up with a formal characterization of state equivalence.





# EQUIVALENT STATES.

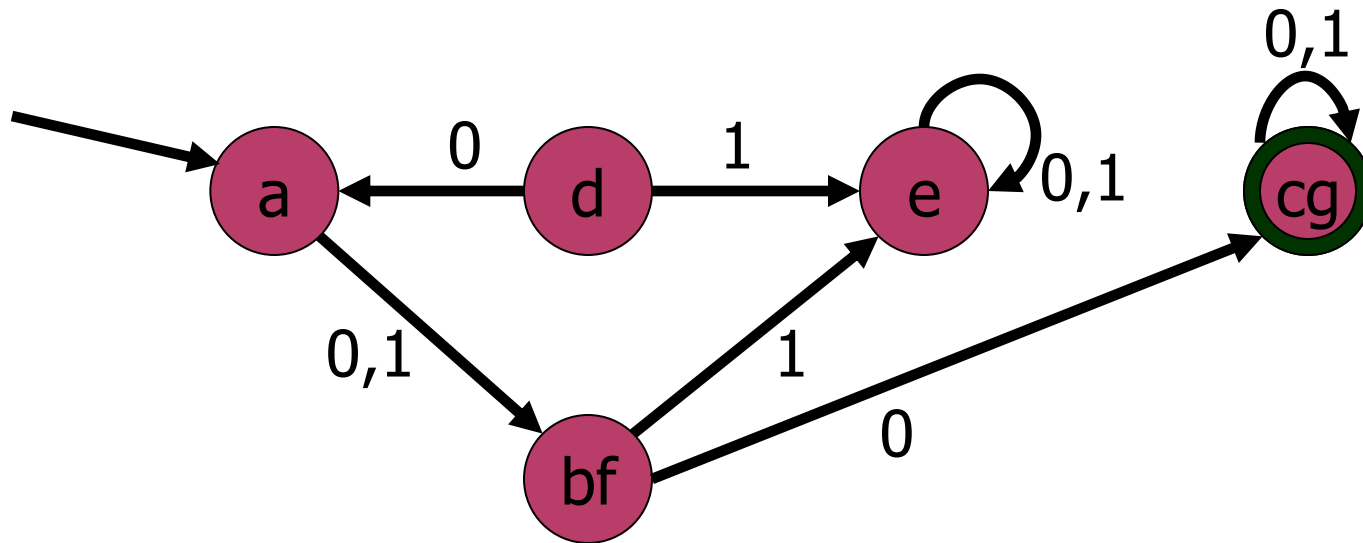
## DEFINITION

DEF: Two states  $q$  and  $q'$  in a DFA  $M = (Q, \Sigma, \delta, q_0, F)$  are said to be ***equivalent*** (or ***indistinguishable***) if for all strings  $u \in \Sigma^*$ , the states on which  $u$  ends on when read from  $q$  and  $q'$  are both accept, or both non-accept.

Equivalent states may be glued together without affecting  $M$ 's behavior.

# FINISHING THE EXAMPLE

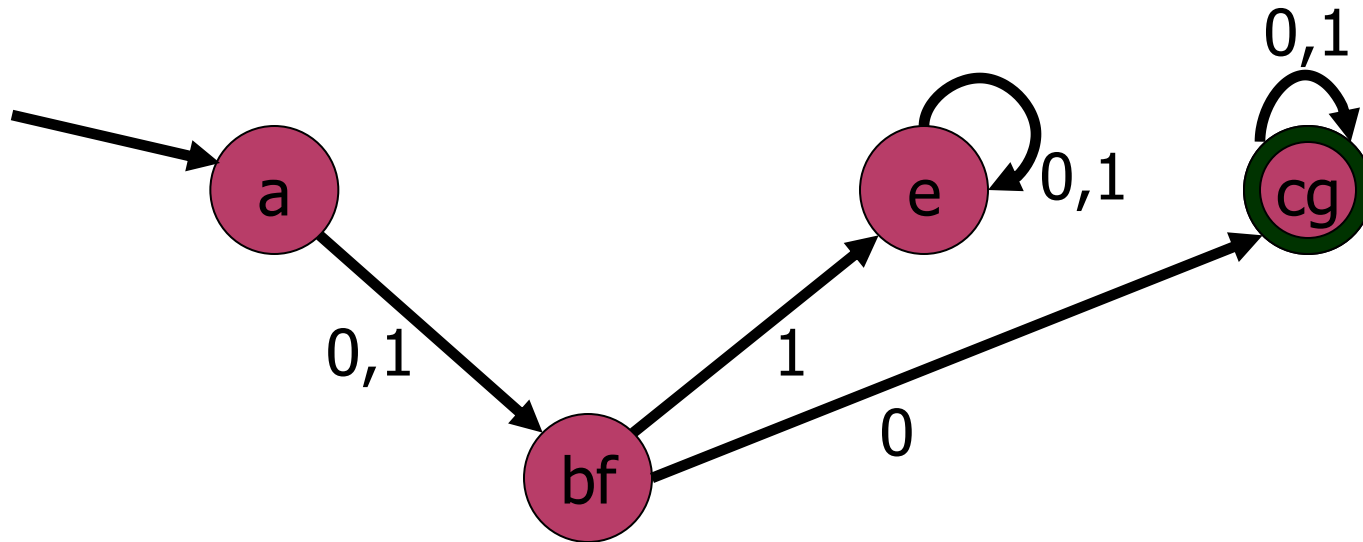
Q: Any other ways to simplify the automaton?



# USELESS STATES

A: Get rid of d.

Getting rid of unreachable *useless states* doesn't affect the accepted language.



# MINIMIZATION ALGORITHM.

## GOALS

DEF: An automaton is *irreducible* if

- it contains no useless states, and
- no two distinct states are equivalent.

The goal of minimization algorithm is to create irreducible automata from arbitrary ones. Later: remarkably, the algorithm actually produces smallest possible DFA for the given language, hence the name “minimization”.

The minimization algorithm *reverses* previous example. Start with least possible number of states, and create new states when forced to.

Explain with a game:

# MINIMIZATION ALGORITHM. (PARTITION REFINEMENT) CODE

```
DFA minimize(DFA (Q,  $\Sigma$ ,  $\delta$ ,  $q_0$ , F) )  
  remove any state  $q$  unreachable from  $q_0$   
  Partition  $P = \{F, Q - F\}$   
  boolean Consistent = false  
  while( Consistent == false )  
    Consistent = true  
    for(every Set  $S \in P$ , char  $a \in \Sigma$ , Set  $T \in P$  )  
      Set temp = { $q \in T \mid \delta(q, a) \in S$  }  
      if (temp !=  $\emptyset$  && temp != T )  
        Consistent = false  
         $P = (P - T) \cup \{temp, T - temp\}$   
  return defineMinimizer( (Q,  $\Sigma$ ,  $\delta$ ,  $q_0$ , F ), P )
```

# MINIMIZATION ALGORITHM. (PARTITION REFINEMENT) CODE

DFA defineMinimizer

(DFA  $(Q, \Sigma, \delta, q_0, F)$ , Partition  $P$ )

Set  $Q' = P$

State  $q'_0 =$  the set in  $P$  which contains  $q_0$

$F' = \{ S \in P \mid S \subseteq F \}$

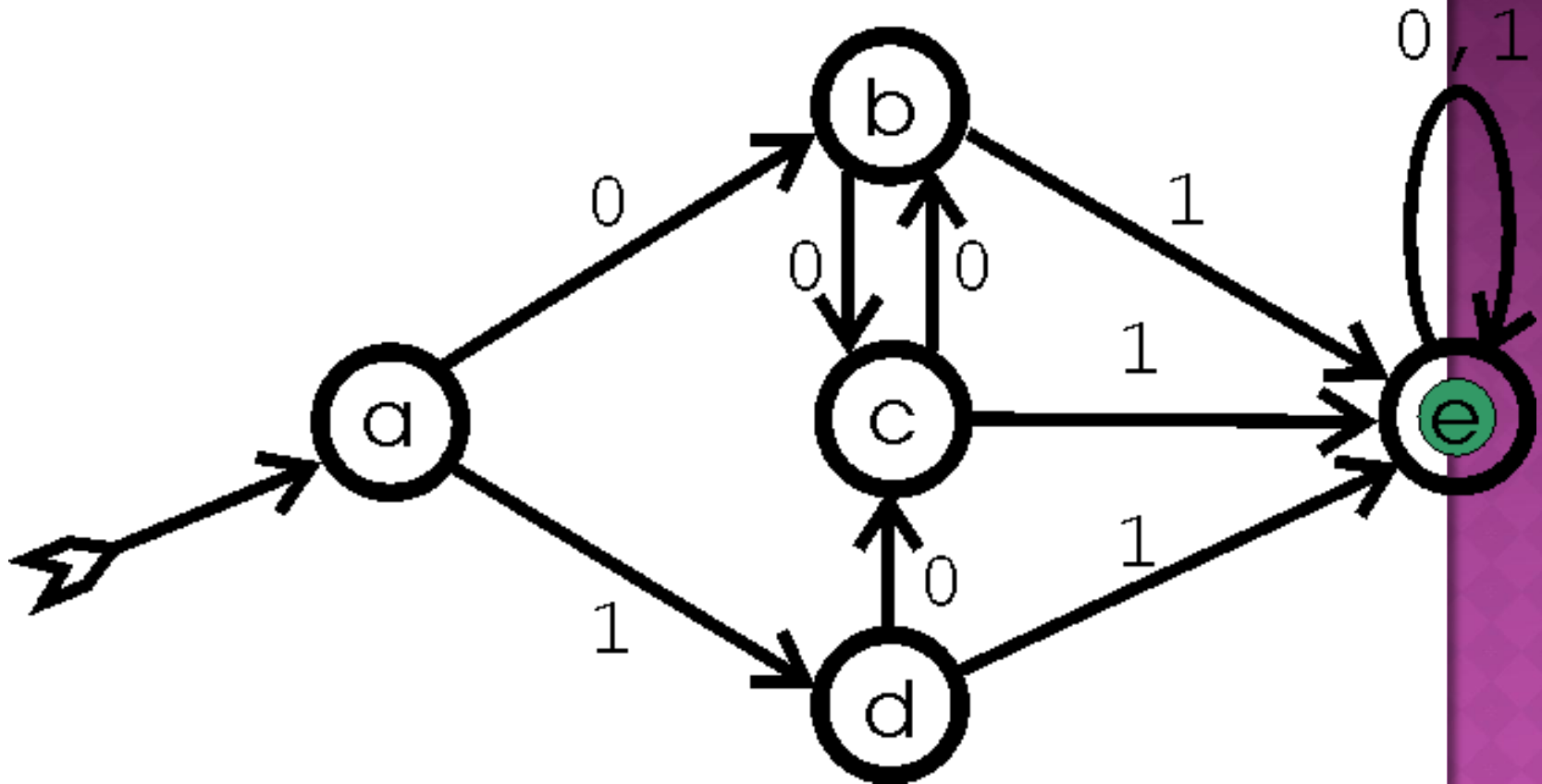
for (each  $S \in P, a \in \Sigma$ )

*define*  $\delta'(S, a) =$  the set  $T \in P$  which contains  
the states  $\delta'(S, a)$

return  $(Q', \Sigma, \delta', q'_0, F')$

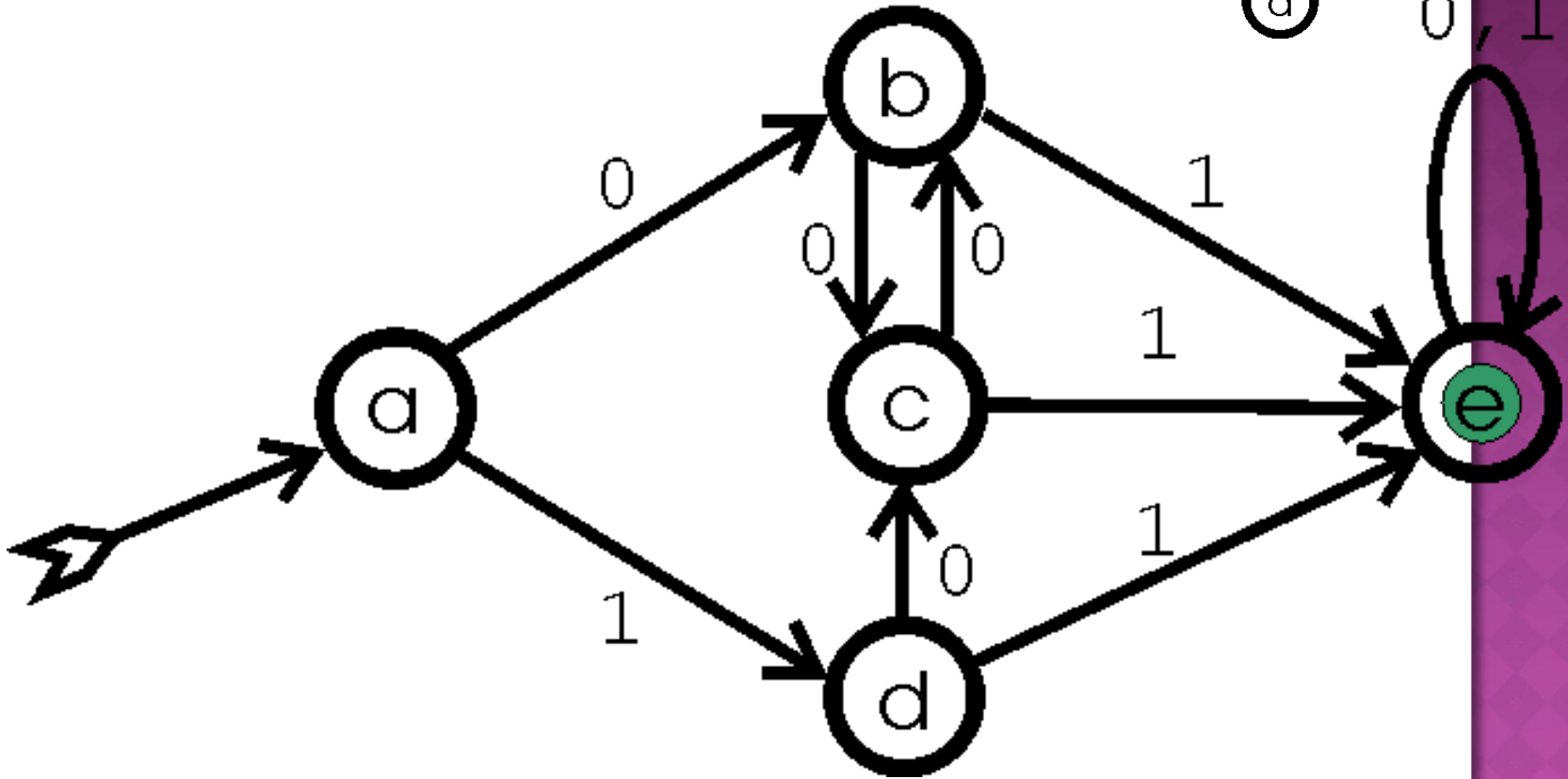
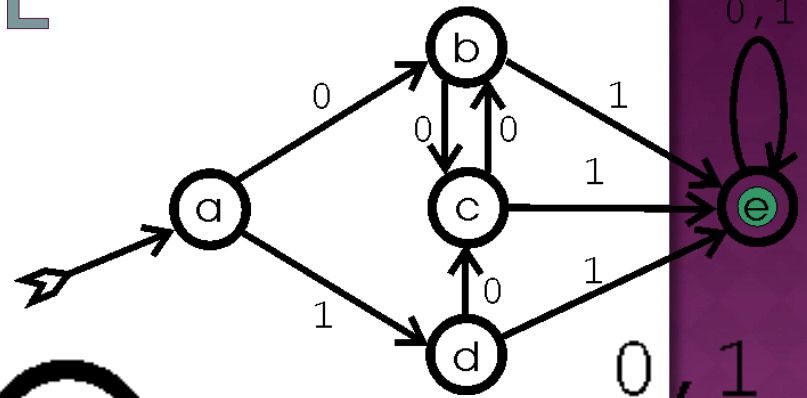
# MINIMIZATION EXAMPLE

Start with a DFA



# MINIMIZATION EXAMPLE

Miniature version →





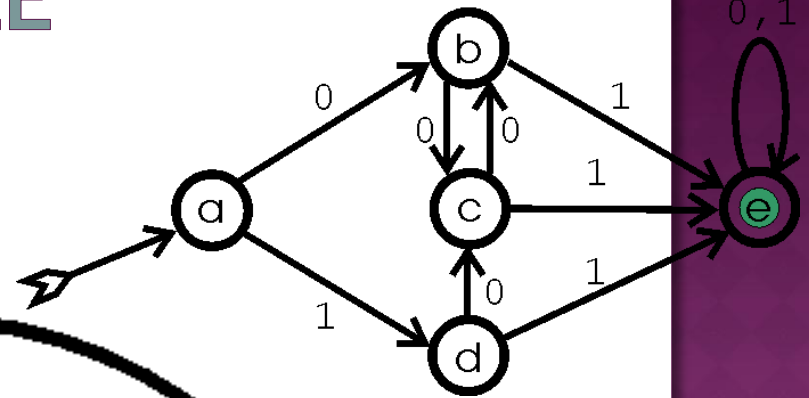
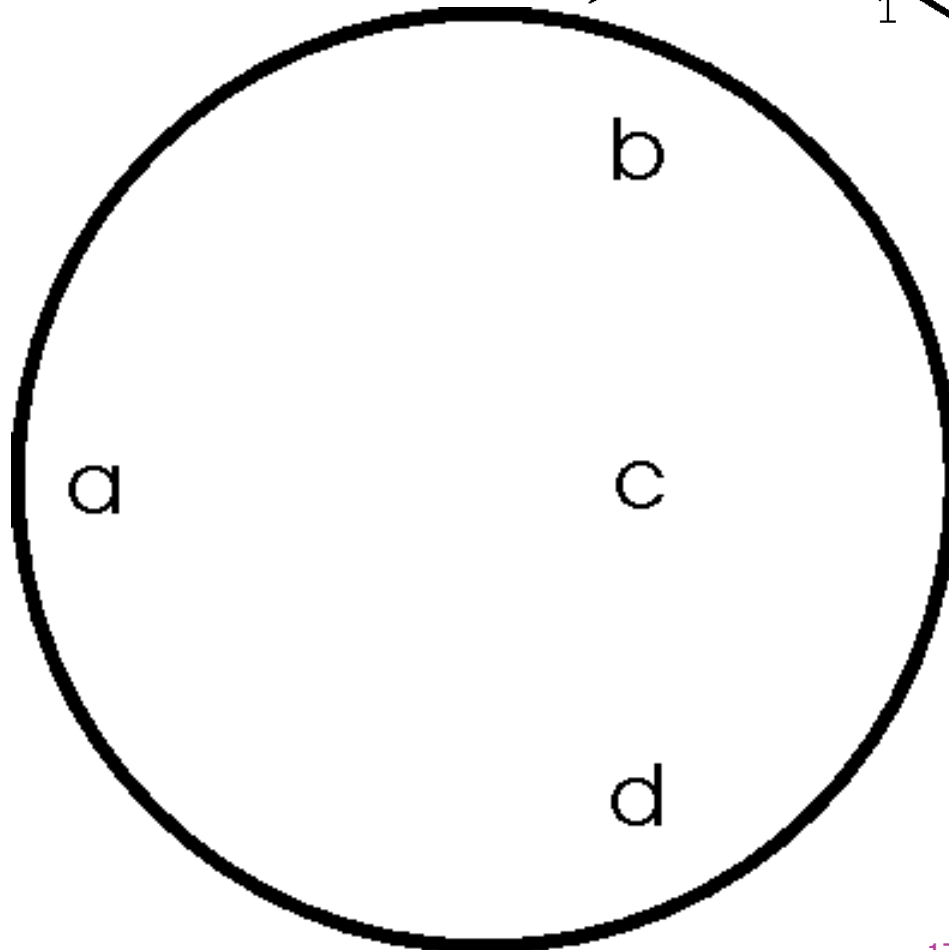
# MINIMIZATION EXAMPLE

Split into two teams.

ACCEPT

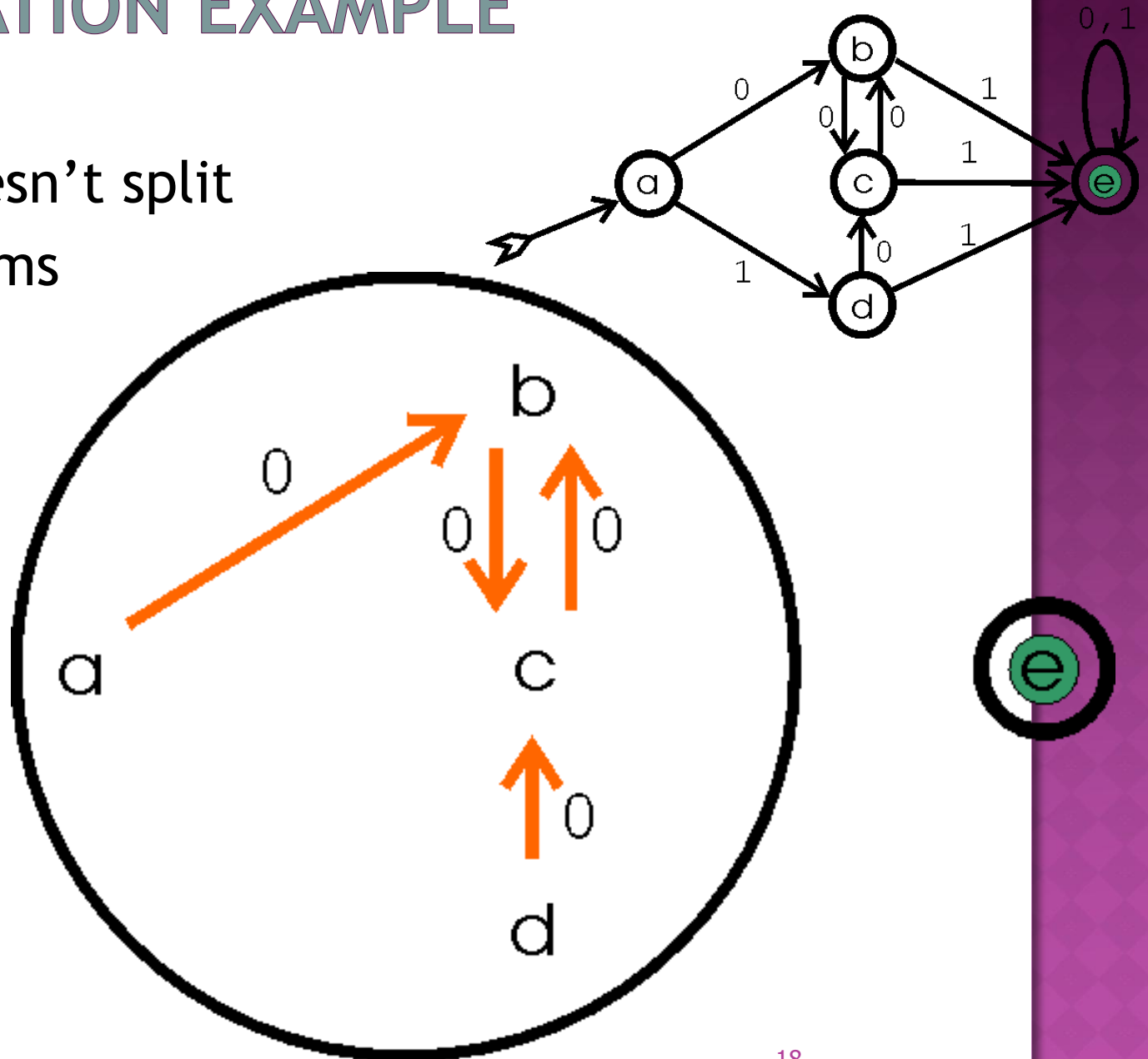
vs.

REJECT



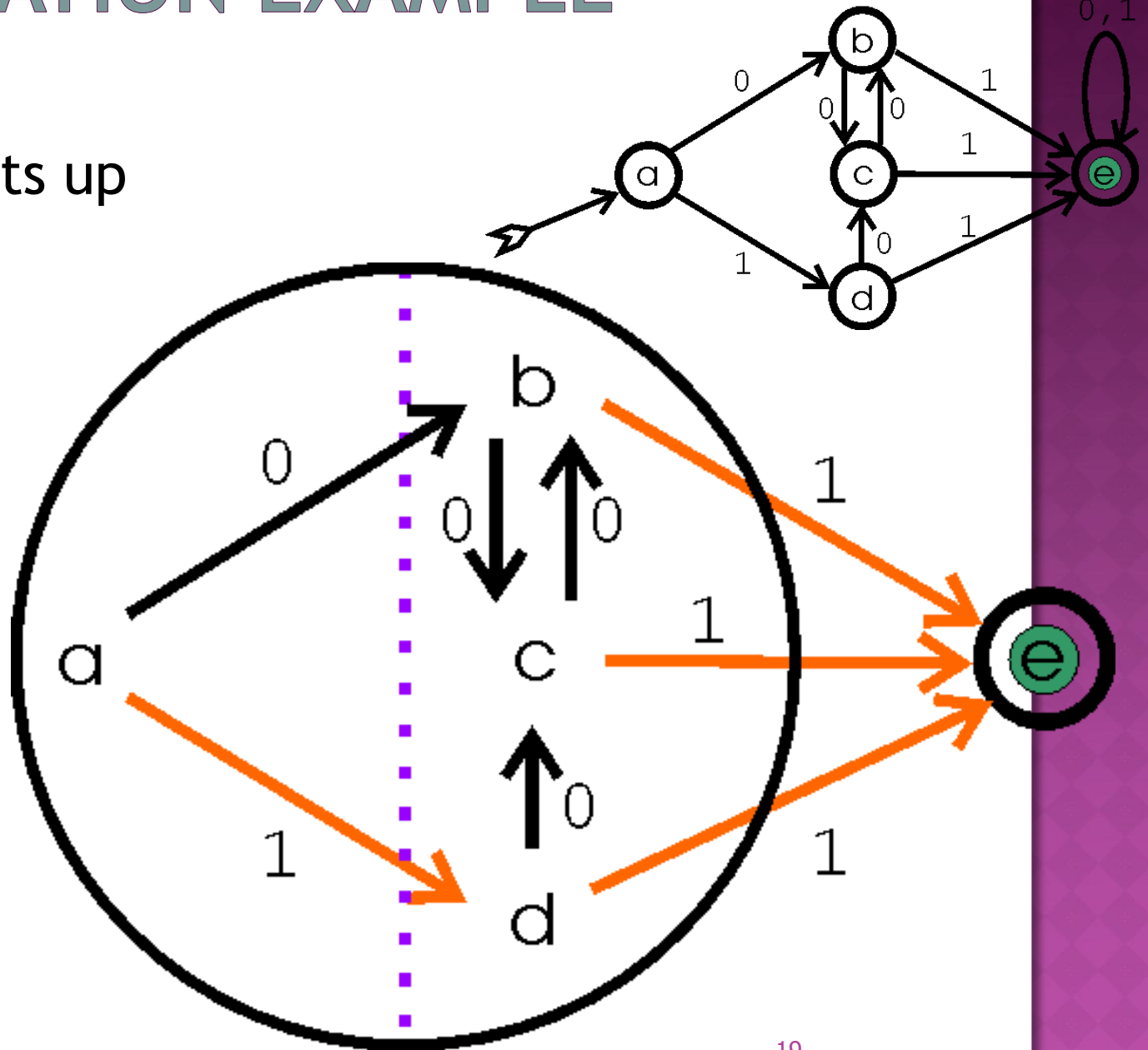
# MINIMIZATION EXAMPLE

0-label doesn't split  
up any teams



# MINIMIZATION EXAMPLE

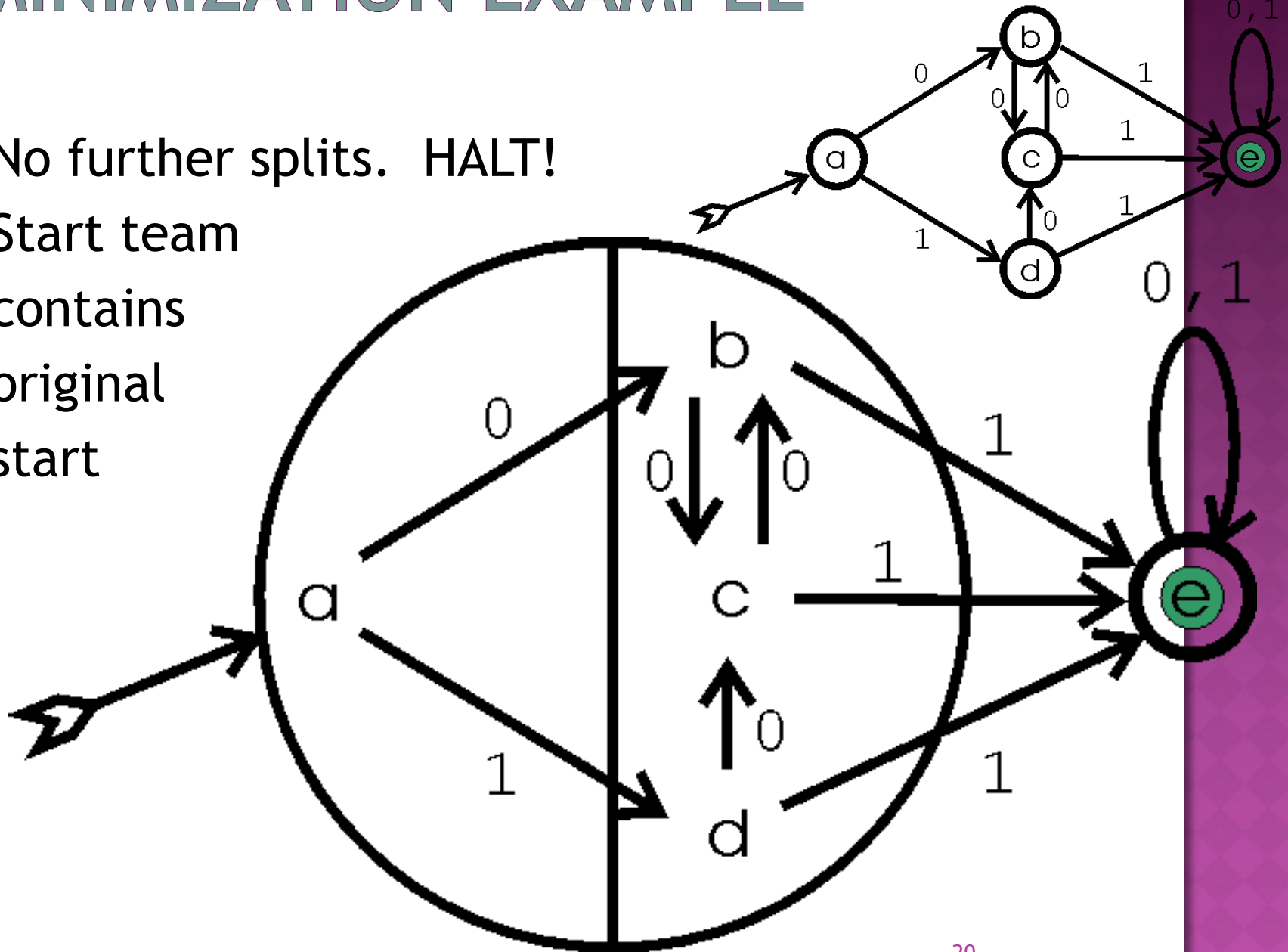
1-label splits up  
REJECT's



# MINIMIZATION EXAMPLE

No further splits. HALT!

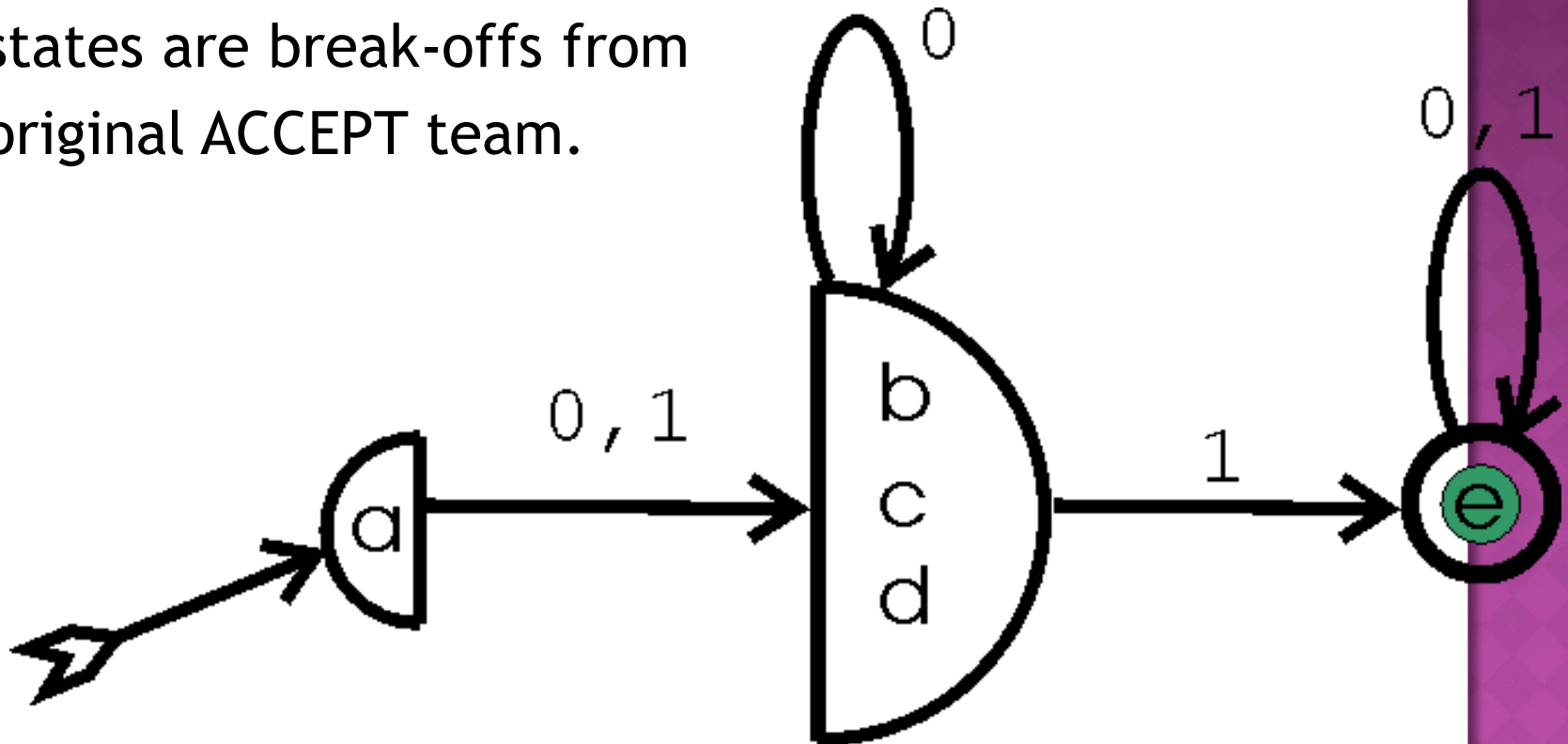
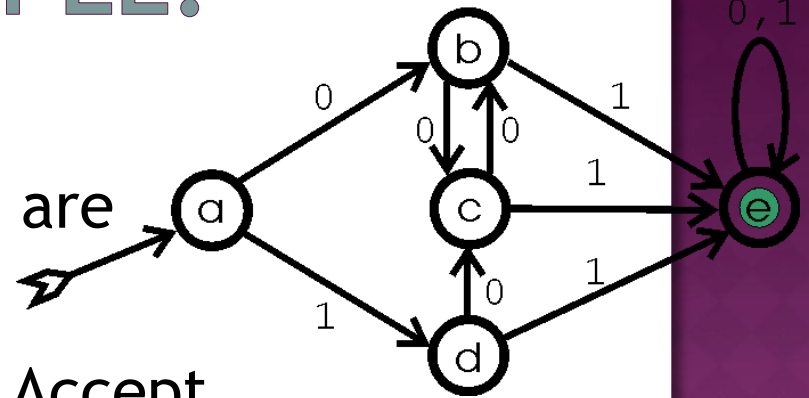
Start team  
contains  
original  
start



# MINIMIZATION EXAMPLE.

## END RESULT

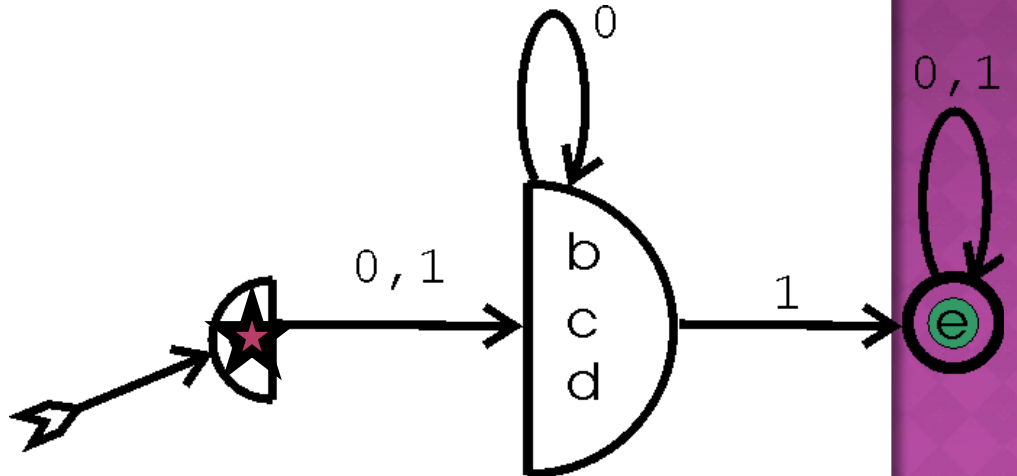
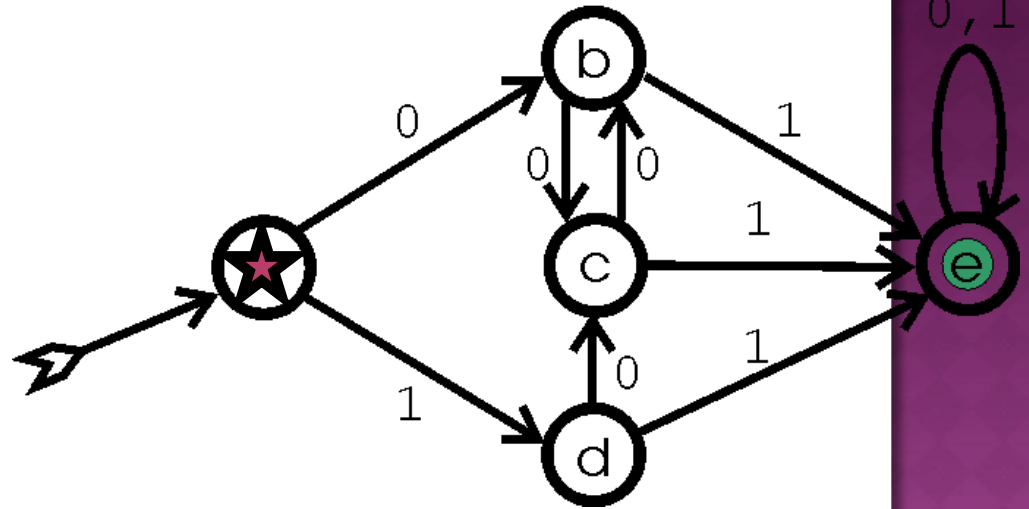
States of the minimal automata are remaining teams. Edges are consolidated across each team. Accept states are break-offs from original ACCEPT team.



# MINIMIZATION EXAMPLE.

## COMPARE

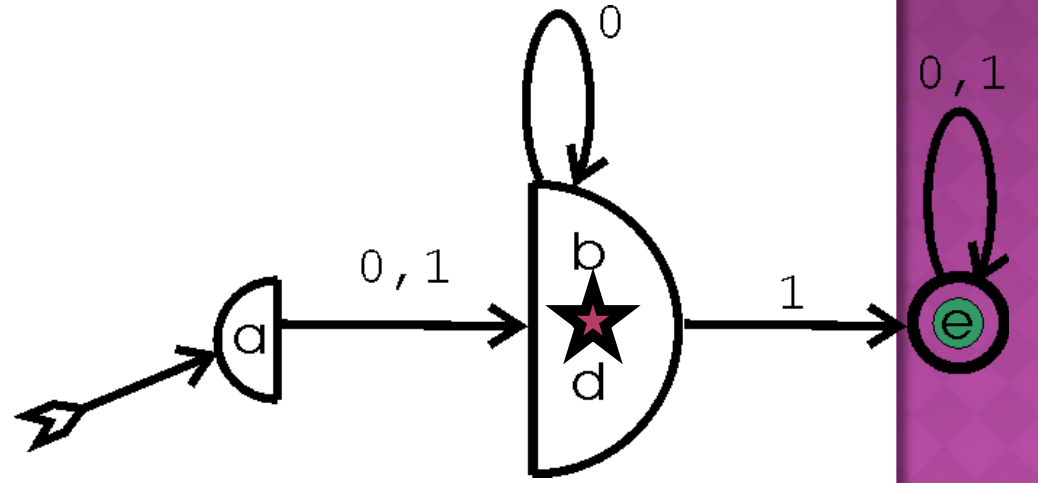
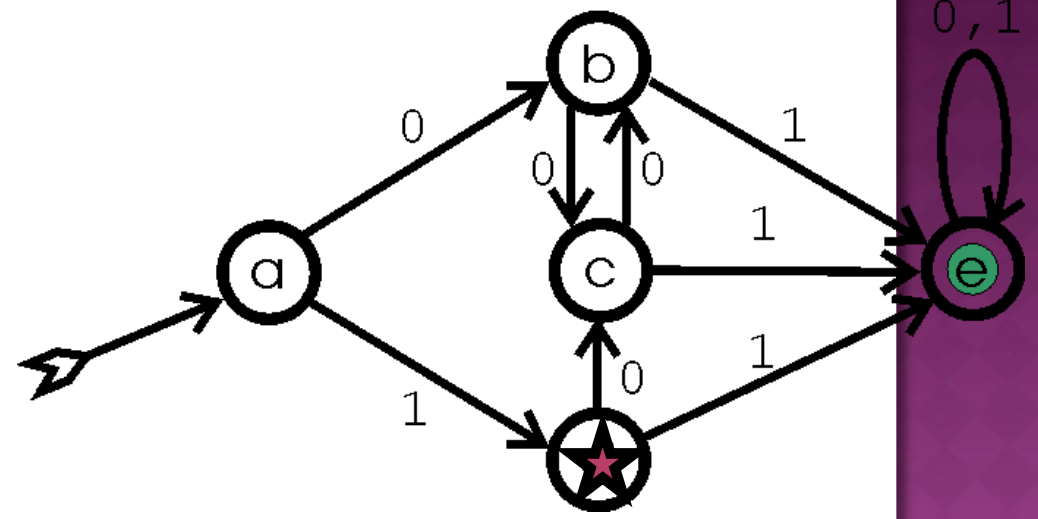
↑  
100100101



# MINIMIZATION EXAMPLE.

## COMPARE

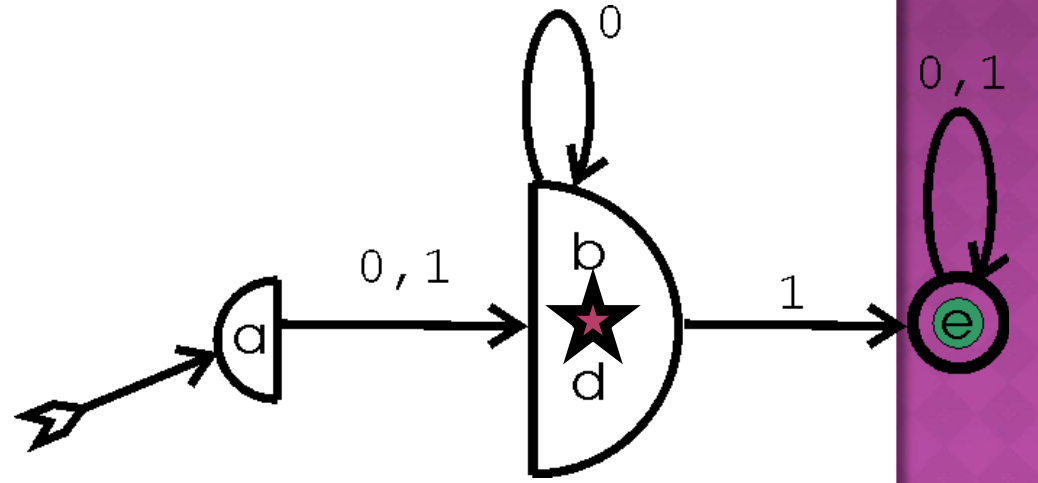
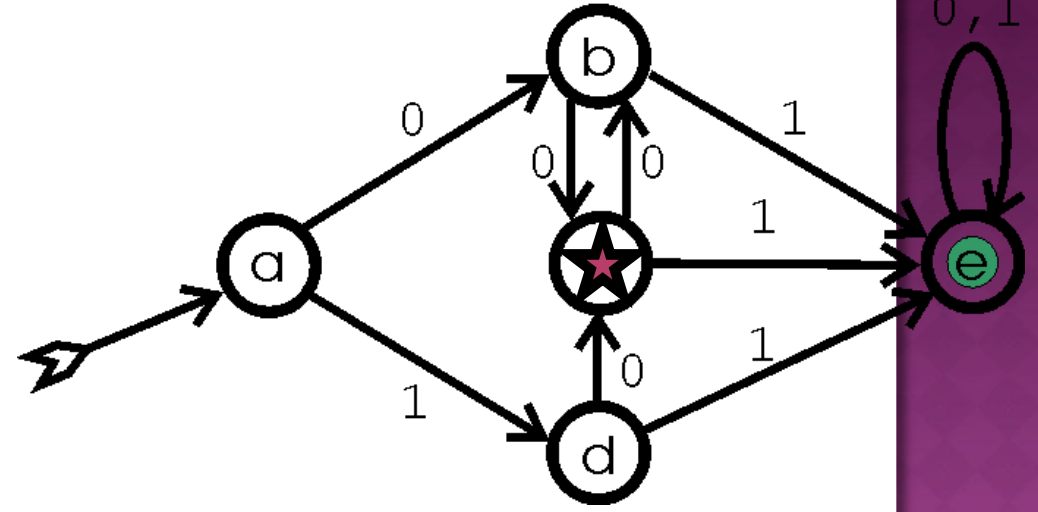
100100101  
↑



# MINIMIZATION EXAMPLE.

## COMPARE

100100101  
↑

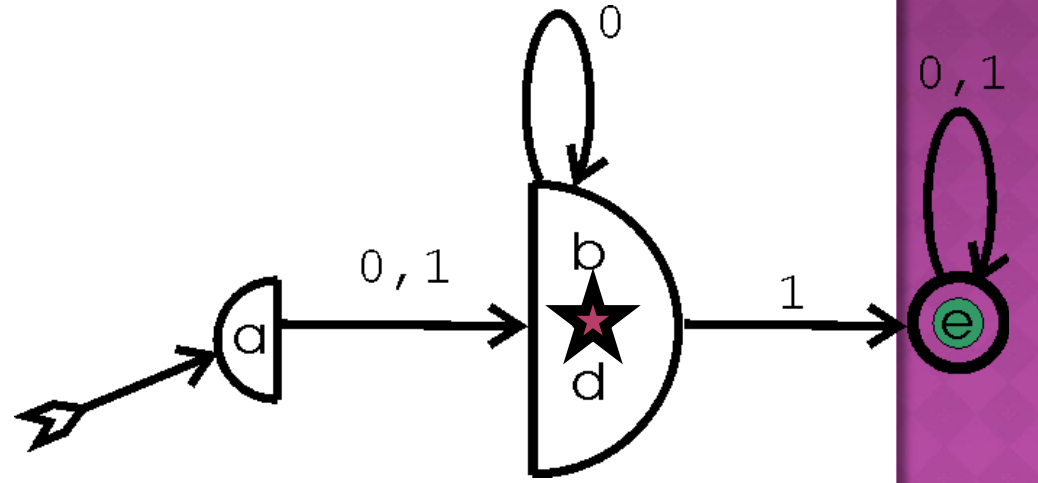
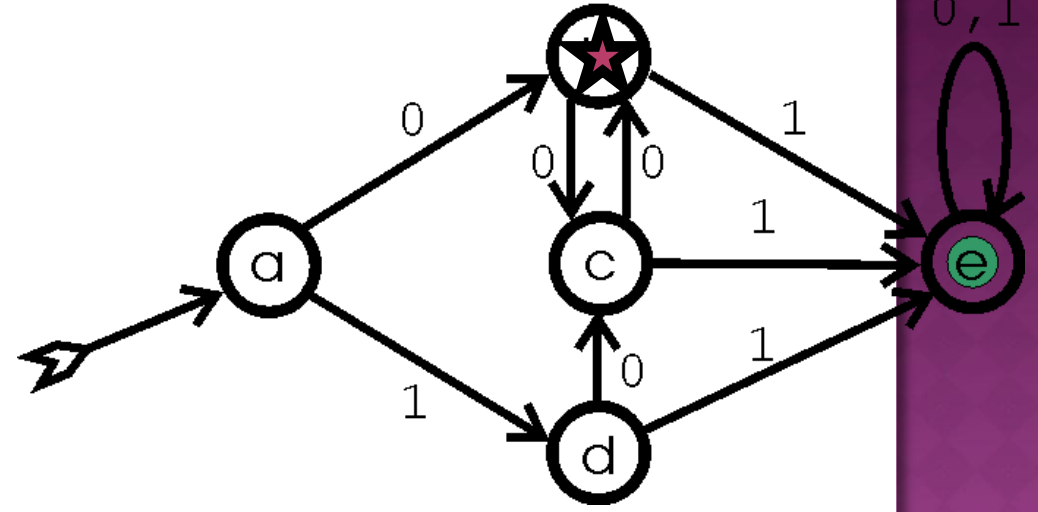




# MINIMIZATION EXAMPLE.

## COMPARE

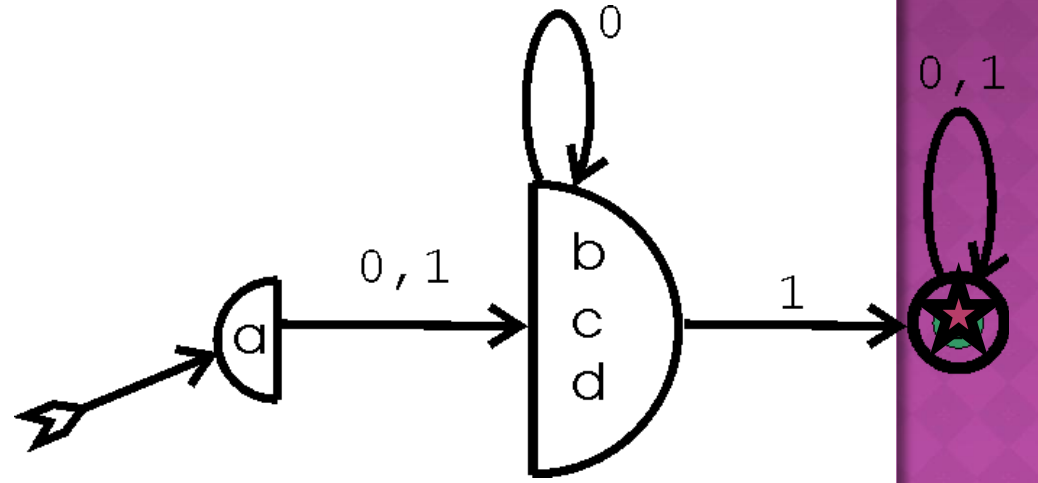
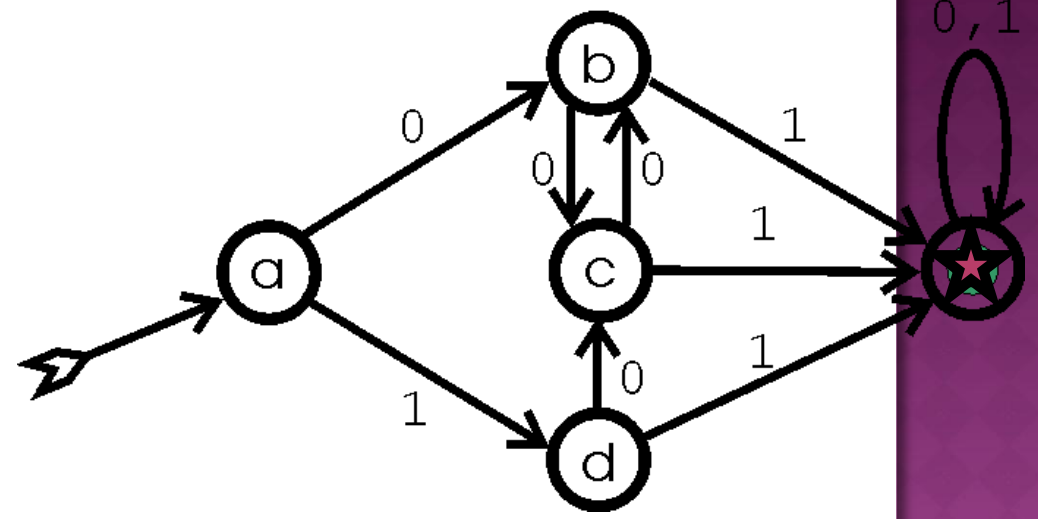
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↑



# MINIMIZATION EXAMPLE.

## COMPARE

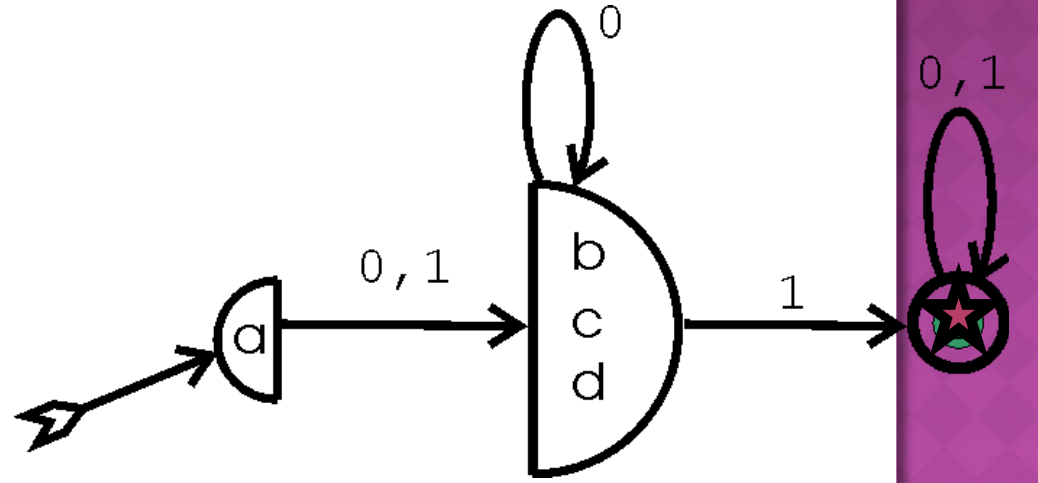
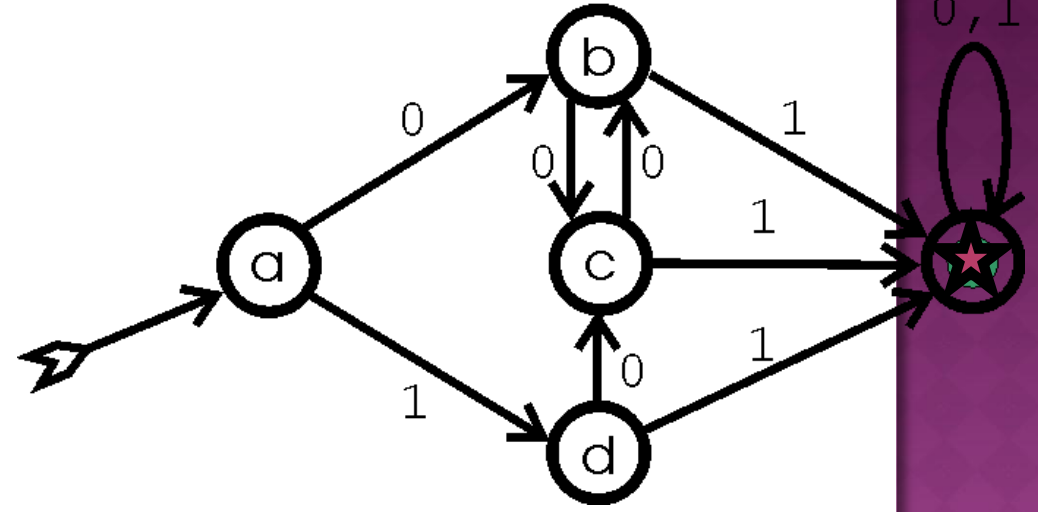
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↑



# MINIMIZATION EXAMPLE.

## COMPARE

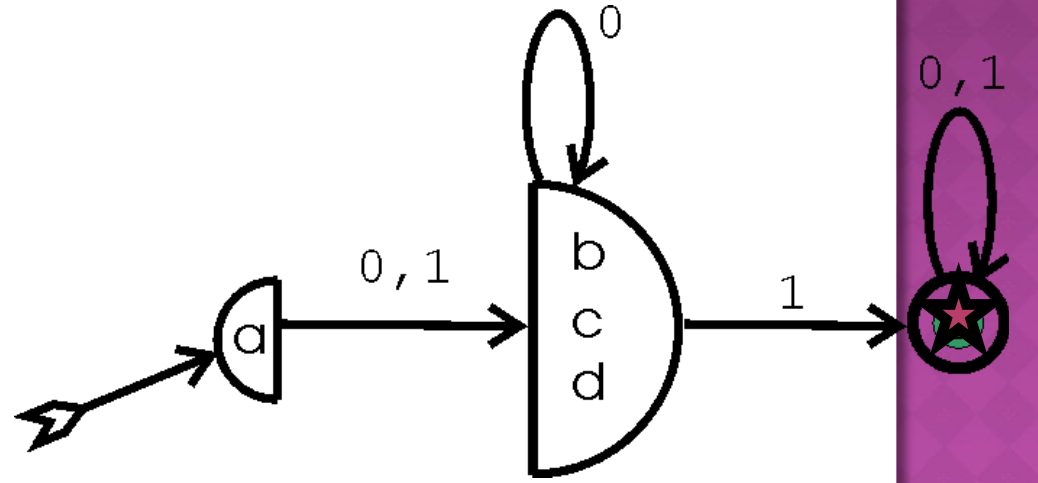
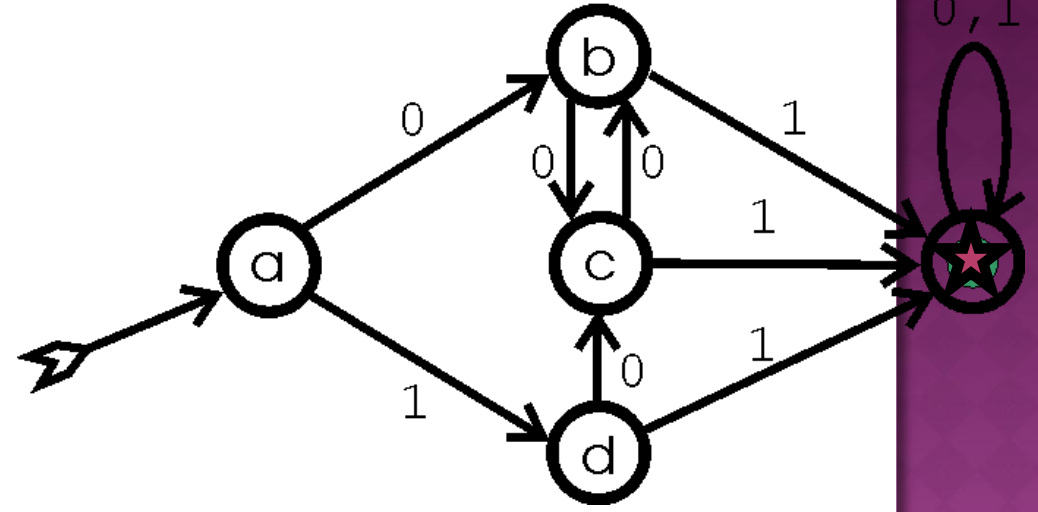
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↑



# MINIMIZATION EXAMPLE.

## COMPARE

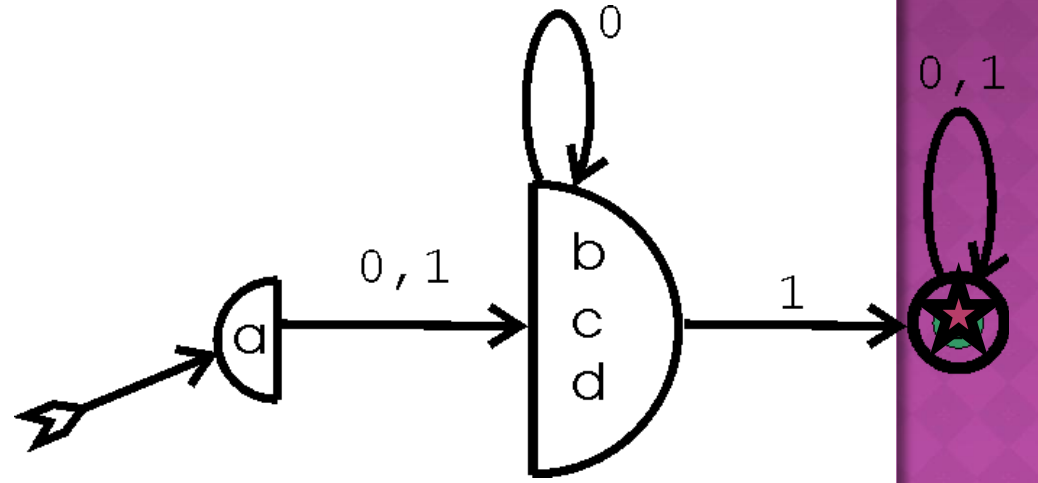
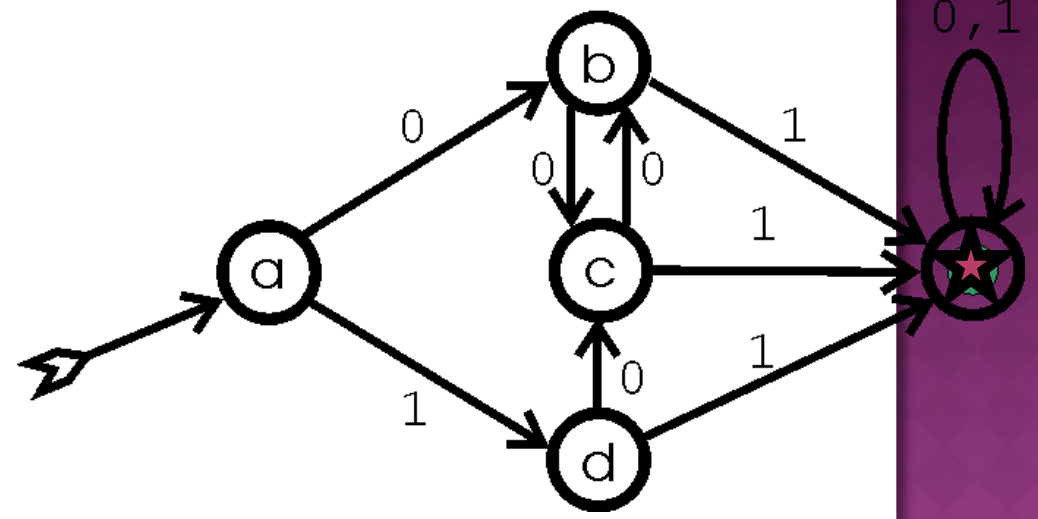
100100101  
↑



# MINIMIZATION EXAMPLE.

## COMPARE

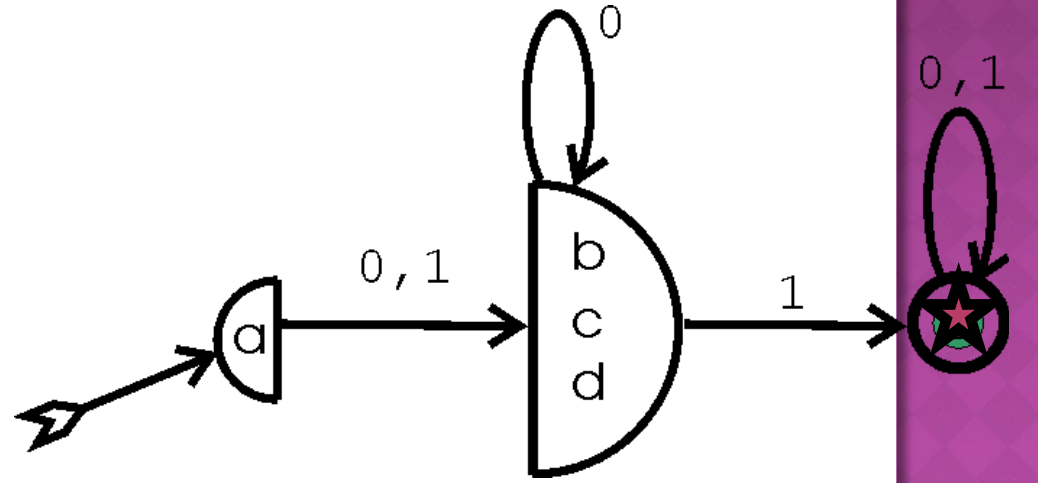
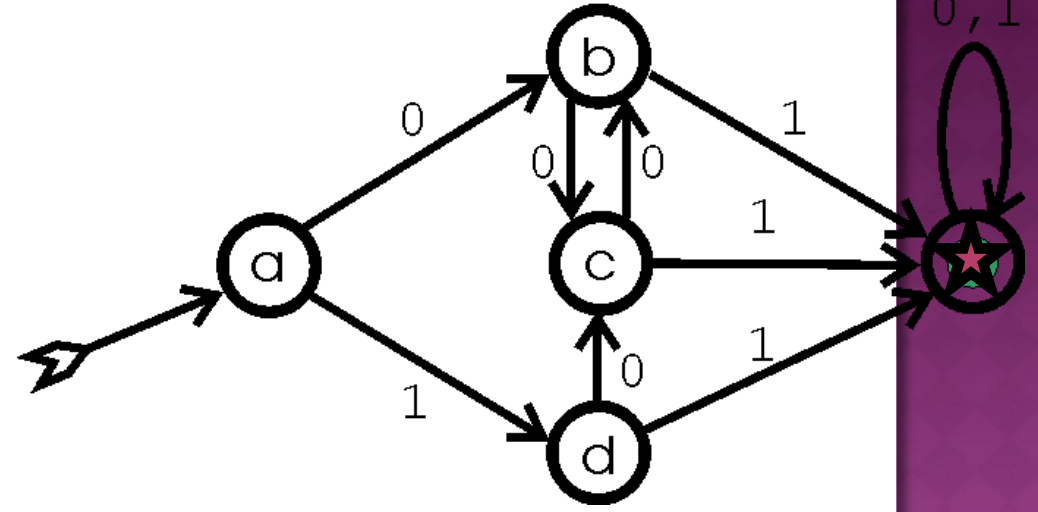
100100101  
↑



# MINIMIZATION EXAMPLE.

## COMPARE

100100101  
↑

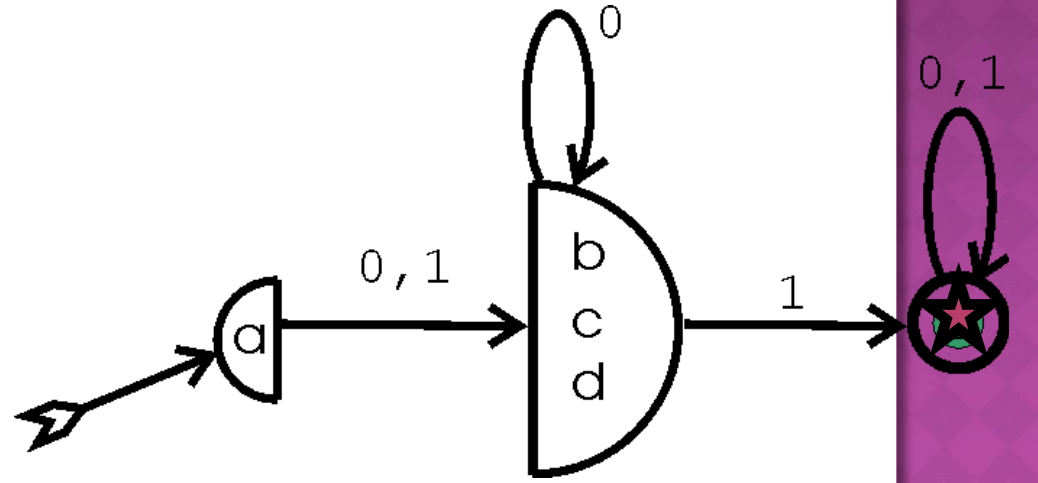
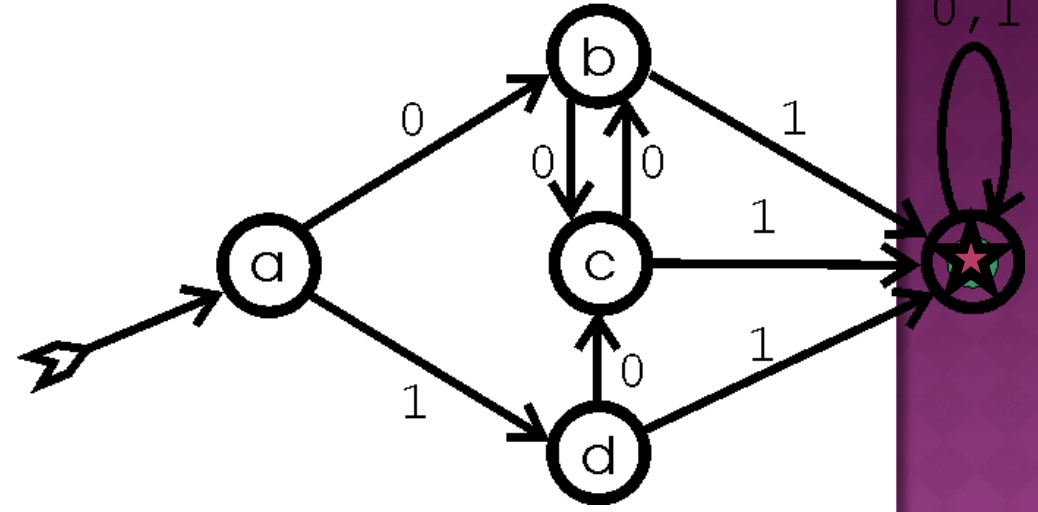


# MINIMIZATION EXAMPLE.

COMPARE

100100101 ↑

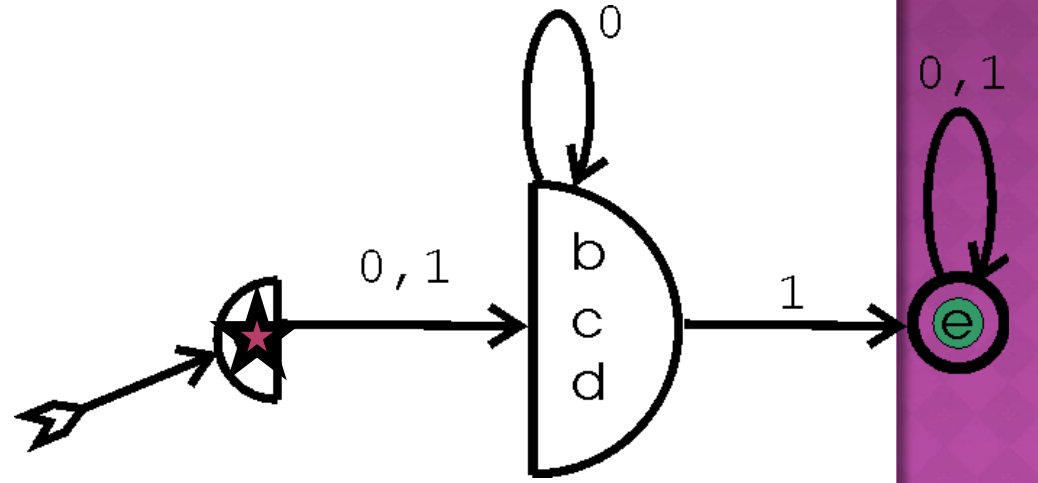
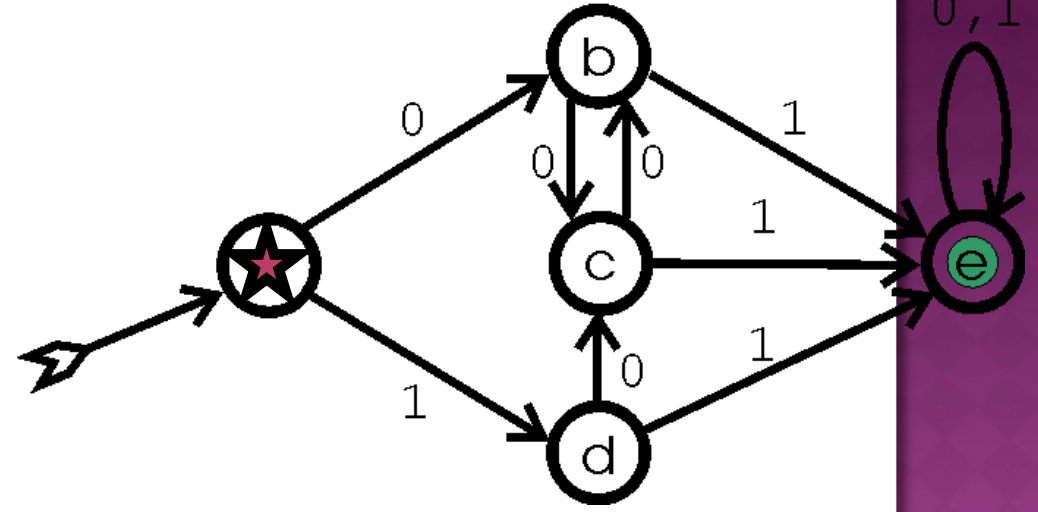
ACCEPTED.



# MINIMIZATION EXAMPLE.

## COMPARE

↑  
10000

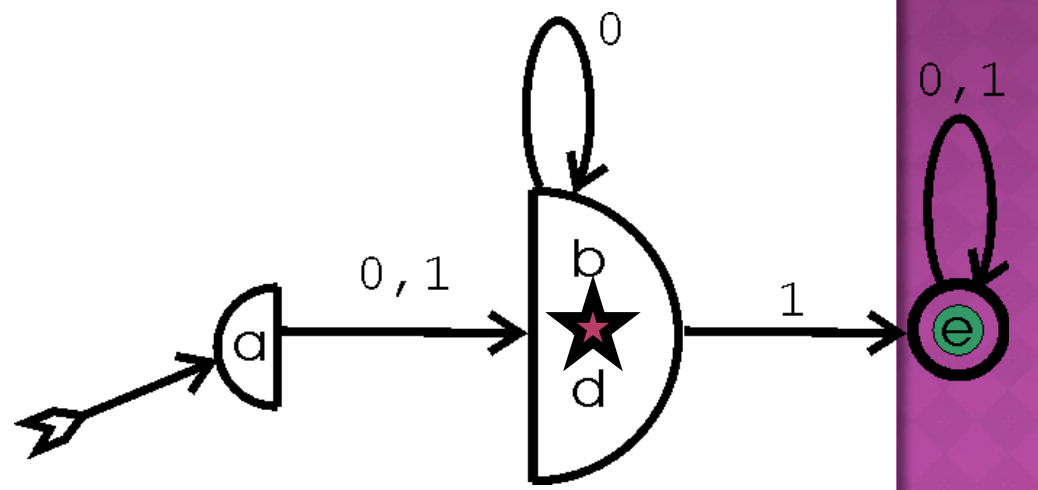
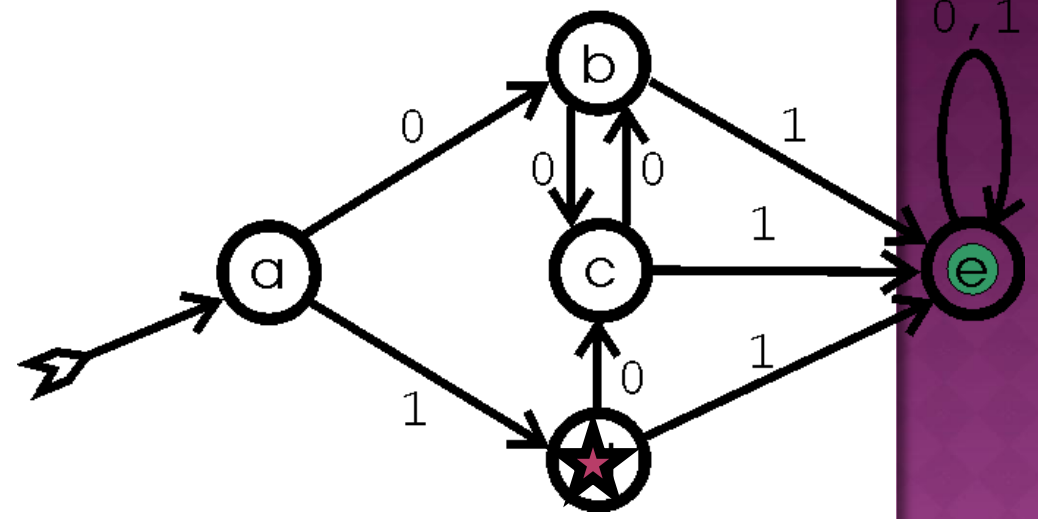




# MINIMIZATION EXAMPLE.

## COMPARE

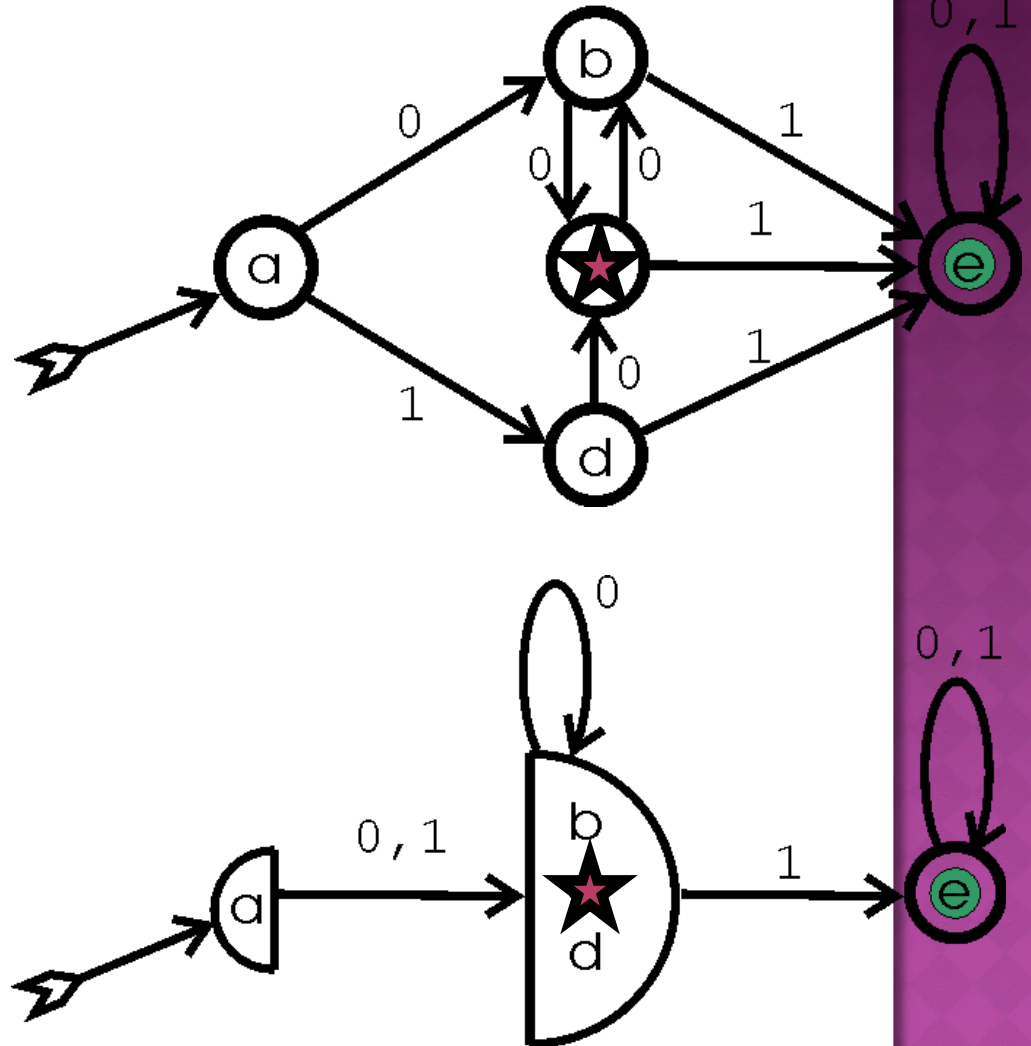
10000  
↑



# MINIMIZATION EXAMPLE.

## COMPARE

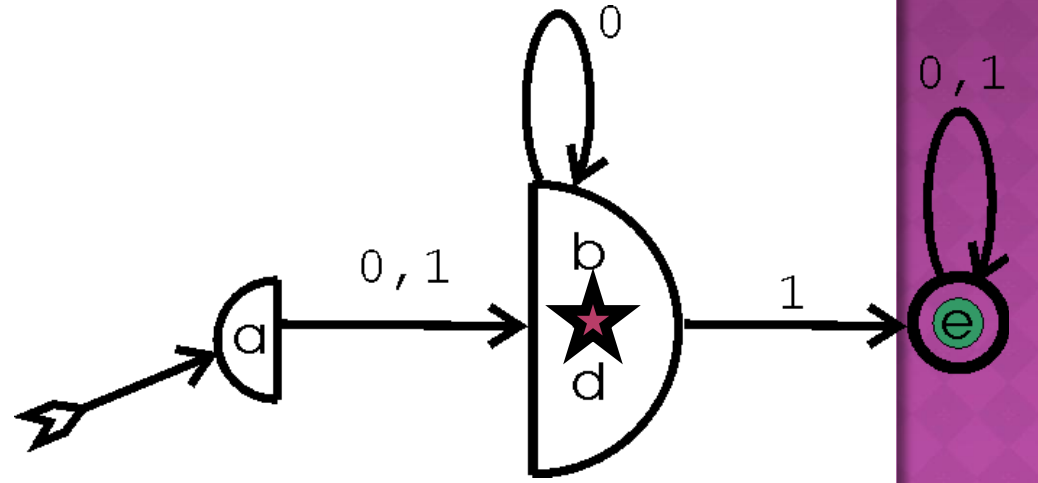
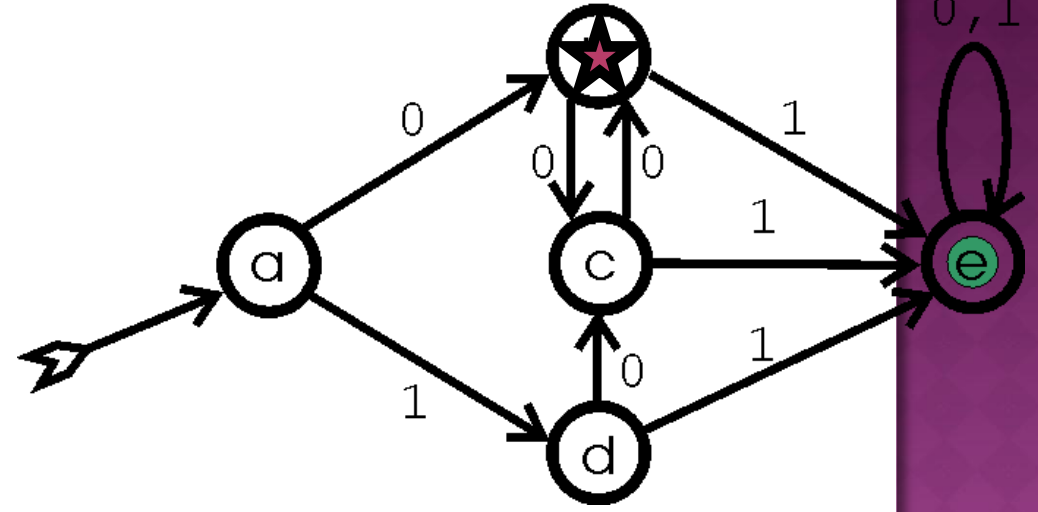
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# MINIMIZATION EXAMPLE.

## COMPARE

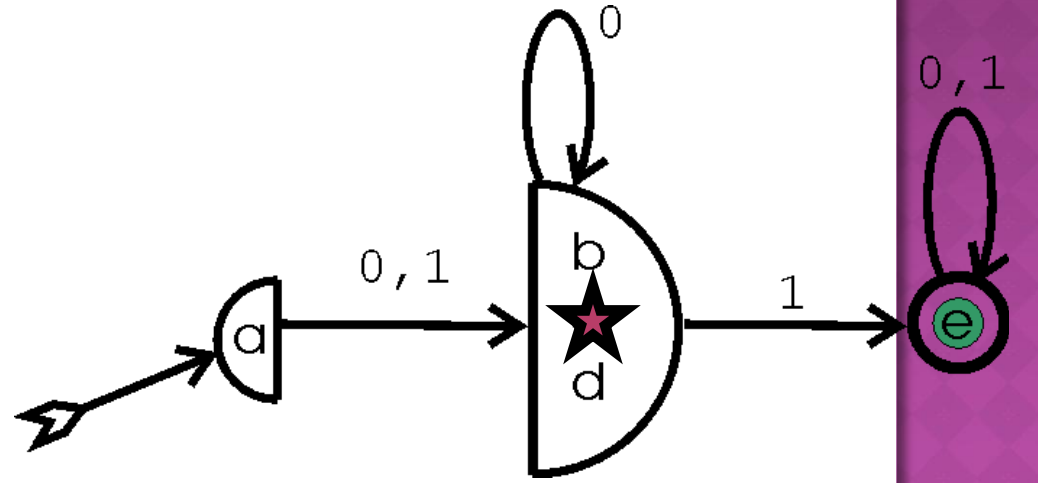
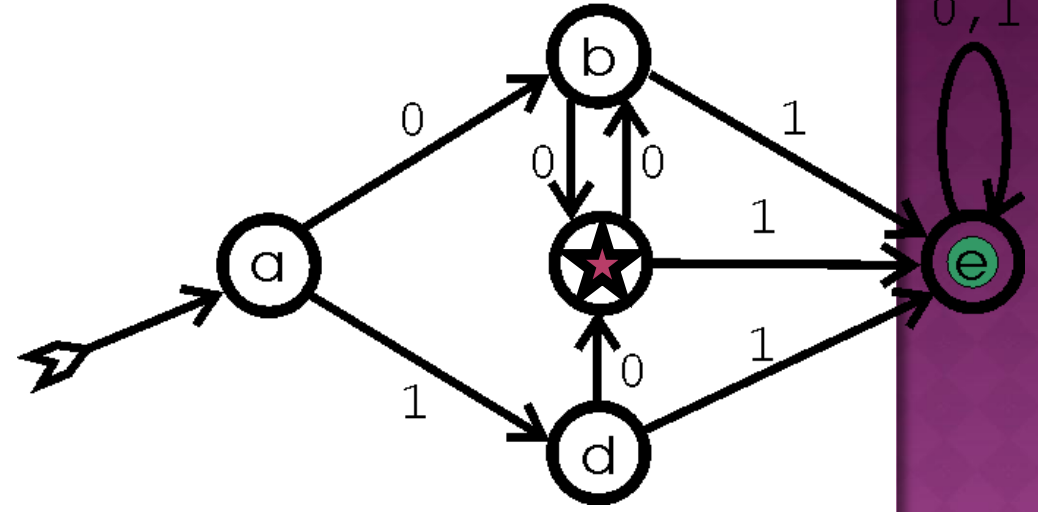
10000  
↑



# MINIMIZATION EXAMPLE.

## COMPARE

10000  
↑



# MINIMIZATION EXAMPLE.

COMPARE

10000  
↑

REJECT.

