

Compiler Design



Lecture-25

The right shift problem

Topics Covered

- The right shift problem
- Addition chains for multiplication

Replace Multiply by Shift

- **$A := A * 4;$**
 - Can be replaced by 2-bit left shift (signed/unsigned)
 - But must worry about overflow if language does
- **$A := A / 4;$**
 - If unsigned, can replace with shift right
 - But shift right arithmetic is a well-known problem
 - Language may allow it anyway (traditional C)

The right shift problem

- Arithmetic Right shift:
 - shift right and use sign bit to fill most significant bits

-5 11111...111111011

SAR 11111...111111101

which is -3, not -2

- in most languages $-5/2 = -2$

Addition chains for multiplication

- If multiply is very slow (or on a machine with no multiply instruction like the original SPARC), decomposing a constant operand into sum of powers of two can be effective:

$$\mathbf{x * 125 = x * 128 - x*4 + x}$$

- two shifts, one subtract and one add, which may be faster than one multiply
- Note similarity with efficient exponentiation method

Folding Jumps to Jumps

- A jump to an unconditional jump can copy the target address

```
JNE lab1
```

```
...
```

```
lab1: JMP lab2
```

Can be replaced by:

```
JNE lab2
```

As a result, lab1 may become dead (unreferenced)

Jump to Return

- A jump to a return can be replaced by a return

```
JMP lab1
```

```
...
```

```
lab1: RET
```

- Can be replaced by

```
RET
```

lab1 may become dead code

Usage of Machine idioms

- Use machine specific hardware instruction which may be less costly.

	$i := i + 1$
ADD $i, \#1$	INC i
	→