

Compiler Design



Lecture-24

Local Optimization



Topics Covered

- Optimization of Basic Blocks
- DAG representation of Basic Block
- Construction of DAG



Optimization of Basic Blocks

- Many structure preserving transformations can be implemented by construction of DAGs of basic blocks

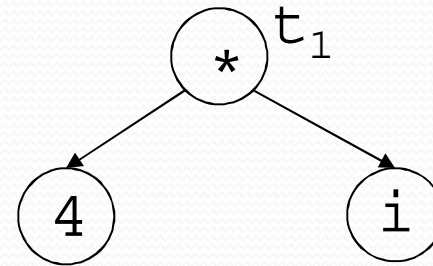


DAG representation of Basic Block (BB)

- Leaves are labeled with unique identifier (var name or const)
- Interior nodes are labeled by an operator symbol
- Nodes optionally have a list of labels (identifiers)
- Edges relates operands to the operator (interior nodes are operator)
- Interior node represents computed value
 - Identifier in the label are deemed to hold the value

Example: DAG for BB

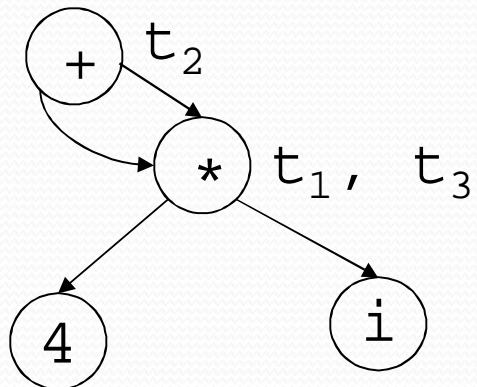
$t_1 := 4 * i$



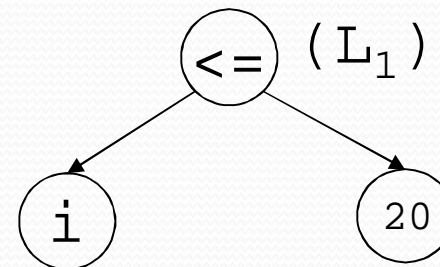
$t_1 := 4 * i$

$t_3 := 4 * i$

$t_2 := t_1 + t_3$



$\text{if } (i \leq 20) \text{ goto } L_1$



Construction of DAGs for BB

- I/p: Basic block, B
- O/p: A DAG for B containing the following information:
 - 1) A label for each node
 - 2) For leaves the labels are ids or consts
 - 3) For interior nodes the labels are operators
 - 4) For each node a list of attached ids (possible empty list, no consts)

Construction of DAGs for BB

- Data structure and functions:
 - Node:
 - 1) Label: label of the node
 - 2) Left: pointer to the left child node
 - 3) Right: pointer to the right child node
 - 4) List: list of additional labels (empty for leaves)
 - **Node (*id*)**: returns the most recent node created for *id*. Else return *undef*
 - **Create(*id,l,r*)**: create a node with label *id* with *l* as left child and *r* as right child. *l* and *r* are optional params.

Construction of DAGs for BB

- Method:

For each 3AC, A in B

A if of the following forms:

1. $x := y \text{ op } z$
2. $x := \text{op } y$
3. $x := y$
1. if $((n_y = \text{node}(y)) == \text{undef})$
 $n_y = \text{Create}(y);$
 if $(A == \text{type } 1)$
 and $((n_z = \text{node}(z)) == \text{undef})$
 $n_z = \text{Create}(z);$

Construction of DAGs for BB

2. If ($A == \text{type 1}$)

Find a node labelled ' op ' with left and right as n_y and n_z respectively [determination of common sub-expression]

If (not found) $n = \text{Create}(op, n_y, n_z);$

If ($A == \text{type 2}$)

Find a node labelled ' op ' with a single child as n_y

If (not found) $n = \text{Create}(op, n_y);$

If ($A == \text{type 3}$) $n = \text{Node}(y);$

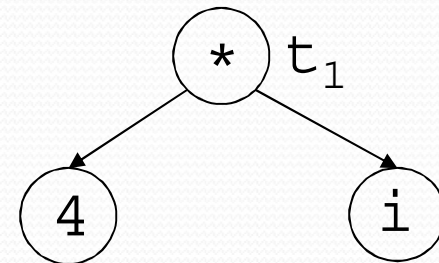
3. Remove x from $\text{Node}(x).\text{list}$

Add x in $n.\text{list}$

$\text{Node}(x) = n;$

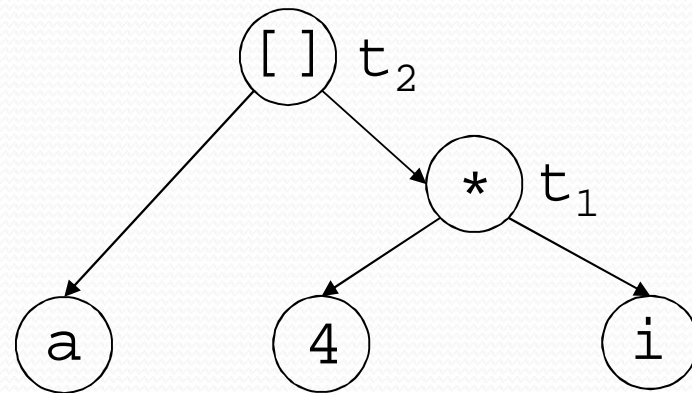
Example: DAG construction from BB

$t_1 := 4 * i$



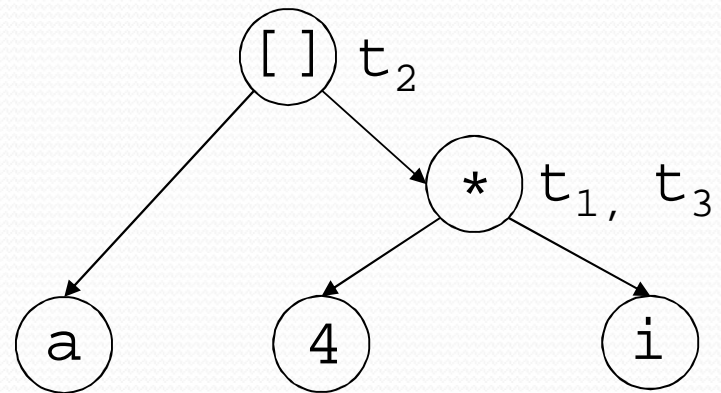
Example: DAG construction from BB

```
t1 := 4 * i  
t2 := a [ t1 ]
```



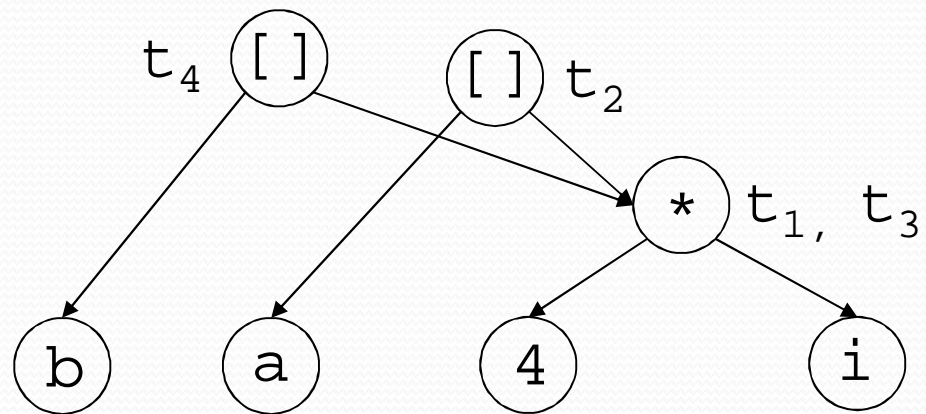
Example: DAG construction from BB

```
t1 := 4 * i  
t2 := a [ t1 ]  
t3 := 4 * i
```



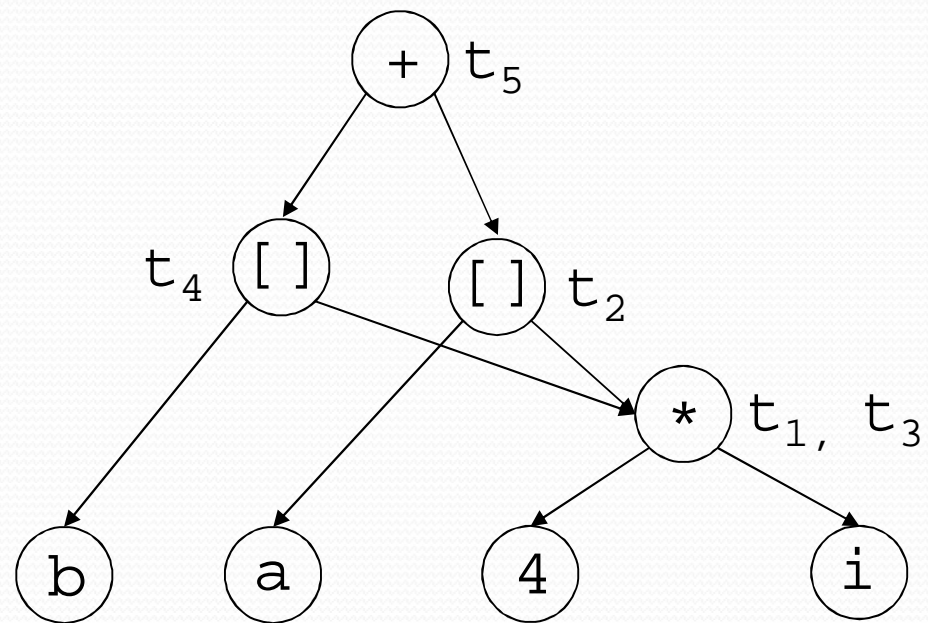
Example: DAG construction from BB

```
t1 := 4 * i  
t2 := a [ t1 ]  
t3 := 4 * i  
t4 := b [ t3 ]
```



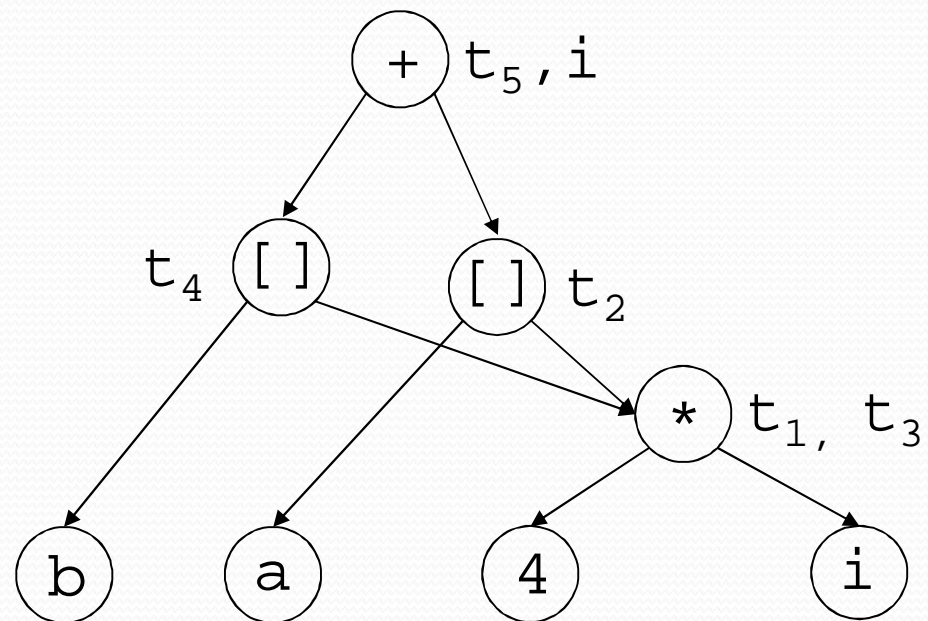
Example: DAG construction from BB

```
t1 := 4 * i  
t2 := a [ t1 ]  
t3 := 4 * i  
t4 := b [ t3 ]  
t5 := t2 + t4
```



Example: DAG construction from BB

```
t1 := 4 * i  
t2 := a [ t1 ]  
t3 := 4 * i  
t4 := b [ t3 ]  
t5 := t2 + t4  
i := t5
```



DAG of a Basic Block

- Observations:
 - A leaf node for the initial value of an id
 - A node n for each statement s
 - The children of node n are the last definition (prior to s) of the operands of n

Optimization of Basic Blocks

- Common sub-expression elimination: by construction of DAG
 - Note: for common sub-expression elimination, we are actually targeting for expressions that compute the same value.

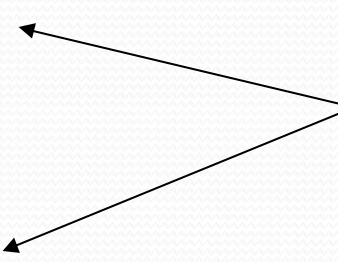
a := b + c

b := b - d

c := c + d

e := b + c

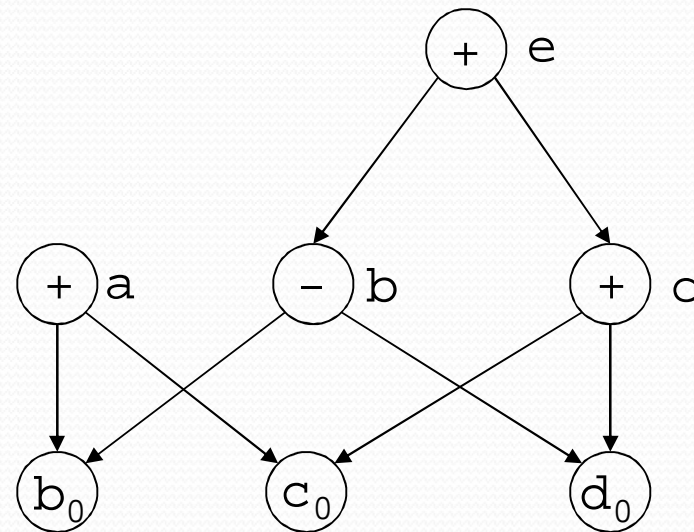
Common expressions
But do not generate the
same result



Optimization of Basic Blocks

- DAG representation identifies expressions that yield the same result

$a := b + c$
$b := b - d$
$c := c + d$
$e := b + c$



Optimization of Basic Blocks

- Dead code elimination: Code generation from DAG eliminates dead code.

