# Lecture 18

## PRINCIPLES OF SATELLITE COMMUNICATION

## **Spreading Codes**

- It is desired that each user's transmitted signal appears noise like and random. Strictly speaking, the signals should appear as Gaussian noise
- Such signals must be constructed from a finite number of randomly preselected stored parameters; to be realizable
- The same signal must be generated at the receiver in perfect synchronization
- We limit complexity by specifying only one bit per sample i.e. a binary sequence

## **Desirable Randomness Properties**

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- Relative frequencies of "o" and "1" should be <sup>1</sup>/<sub>2</sub> (Balance property)
- Run lengths of zeros and ones should be (Run property):
  - Half of all run lengths should be unity
  - One quarter should be of length two
  - One eighth should be of length three
  - A fraction 1/2<sup>n</sup> of all run lengths should be of length n for all finite n

#### Desirable Randomness Properties (contd...)

• If the random sequence is shifted by any nonzero number of elements, the resulting sequence should have an equal number of agreements and disagreements with the original sequence



- A deterministically generated sequence that nearly satisfies these properties is referred to as a Pseudorandom Sequence (PN)
- Periodic binary sequences can be conveniently generated using linear feedback shift registers (LFSR)
- If the number of stages in the LFSR is r,  $P \leq 2^r$  1 where P is the period of the sequence

#### PN Sequences (contd...)

However, if the feedback connections satisfy a specific property, P = 2<sup>r</sup> - 1. Then the sequence is called a Maximal Length Shift Register (MLSR) or a PN sequence.

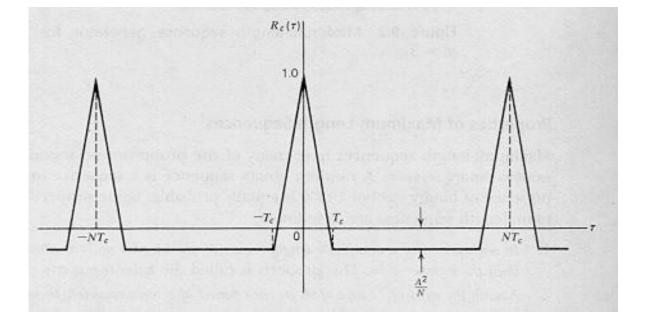
• Thus if r=15, P=32767.

## **Randomness Properties of PN Sequences**

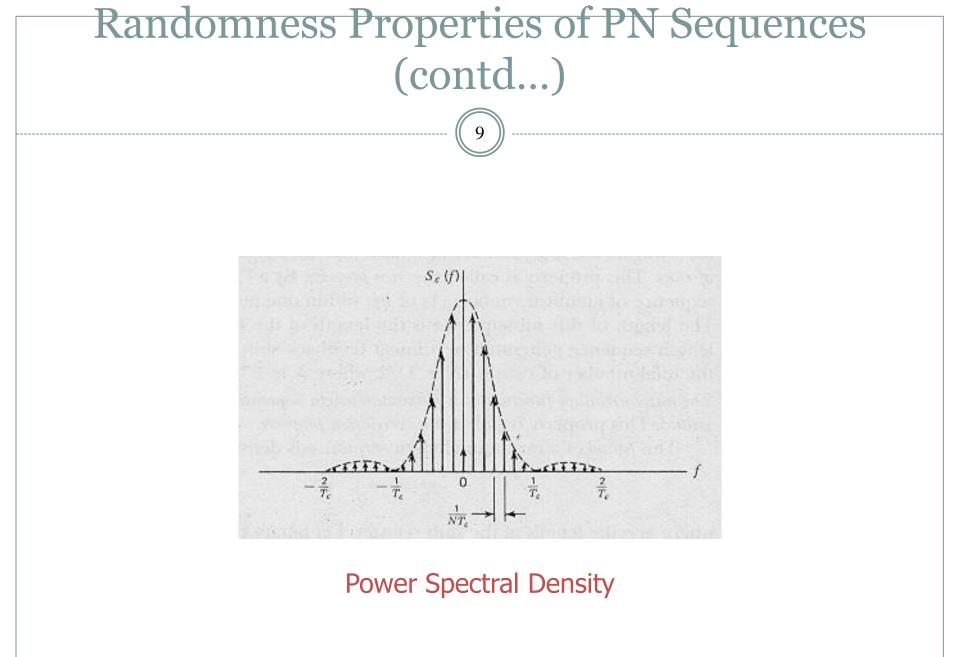
- Balance property Of the 2<sup>r</sup> 1 terms, 2<sup>r-1</sup> are one and 2<sup>r-1</sup>-1 are zero. Thus the unbalance is 1/P. For r=50; 1/P≅10<sup>-15</sup>
- Run property Relative frequency of run length n (zero or ones) is 1/  $2^n$  for  $n \le r-1$  and  $1/(2^r 1)$  for n = r
- One run length each of r-1 zeros and r ones occurs. There are no run lengths for n > r
- Autocorrelation property The number of disagreements exceeds the number of agreements by unity. Thus again the discrepancy is 1/p

## Randomness Properties of PN Sequences (contd.)

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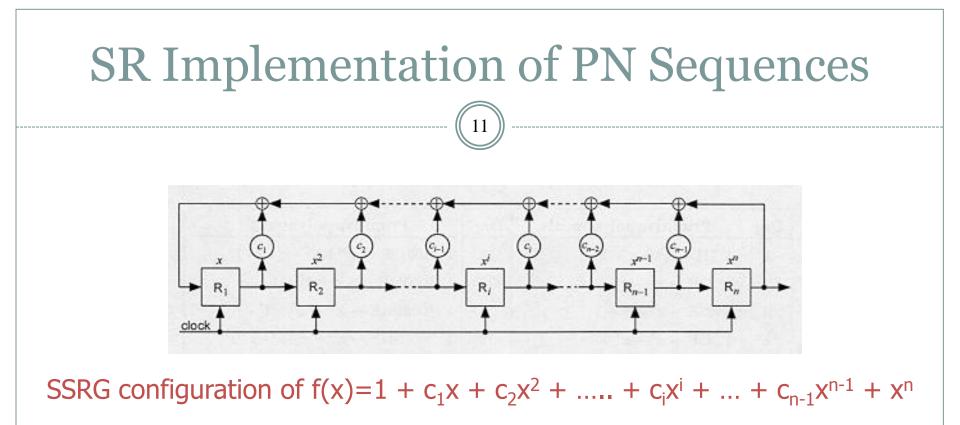
Autocorrelation function

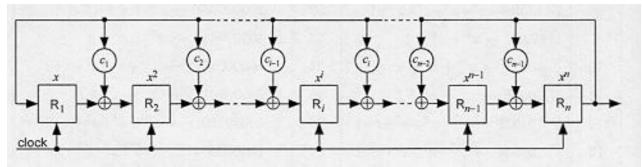


## SR Implementation of PN Sequences

- The feedback connection should correspond to a primitive polynomial.
- Primitive polynomials of every degree exist. The number of primitive polynomials of degree r is given by :

- Simple Shift Register Generator (SSRG) Fibonacci configuration.
- Modular Shift Register Generator (MSRG) Galois configuration.





MSRG configuration of  $f(x)=1 + c_1x + c_2x^2 + \dots + c_ix^i + \dots + c_{n-1}x^{n-1} + x^n$ 

## **PN Sequences Specified in IS-95**

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- A **long** PN sequence (r =42) is used to scramble the user data with a different code shift for each user
- The 42-degree characteristic polynomial is given by:
  - $\stackrel{\text{O}}{\times} x^{42} + x^{41} + x^{40} + x^{39} + x^{37} + x^{36} + x^{35} + x^{32} + x^{26} + x^{25} + x^{24} + x^{23} + x^{21} + x^{20} + x^{17} + x^{16} + x^{15} + x^{11} + x^9 + x^7 + 1$

• The period of the long code is  $2^{42} - 1 \approx 4.4^{*10^2}$  chips and lasts over 41 days

### PN Sequences Specified in IS-95 (contd...)

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- Two "short" PN sequences (r=15) are used to spread the quadrature components of the forward and reverse link waveforms
- The characteristic polynomials are given by : •  $x^{15}+x^{10}+x^8+x^7+x^6+x^2+x$  (I-channel) •  $x^{15}+x^{12}+x^{11}+x^{10}+x^9+x^5+x^4+x^3+1$  (Q-channel)
- The period of the short code is:  $2^{15} - 1 = 32767$  chips  $\equiv 80/3$  ms

## Orthogonal Spreading Codes – Walsh Codes

Walsh functions of order N are defined as a set of N time functions denoted as  $\{W_i(t); t \in (0,T), j=0,1,...N-1\}$  such that:

- W<sub>j</sub>(t) takes on the values {+1, -1} except at the jumps, where it takes the value zero
- $W_j(t) = 1$  for all j
- $W_{j}(t)$  has precisely j sign changes in the interval (0,T)

- Each  $W_j(t)$  is either even or odd with respect to T/2 i.e. the mid point