A decorative vertical bar on the left side of the slide. It consists of a dark teal background with a white vertical stripe. To the right of the teal bar are several orange circles of varying sizes, arranged in a cluster. The title text is positioned to the right of this decorative bar.

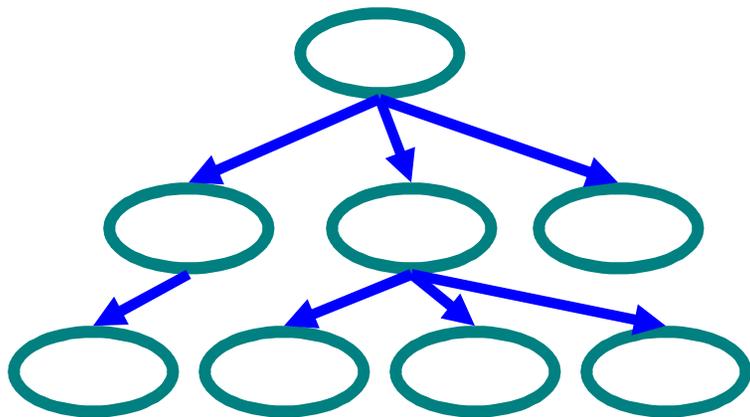
# DATA STRUCTURES USING 'C'

# Lecture-23

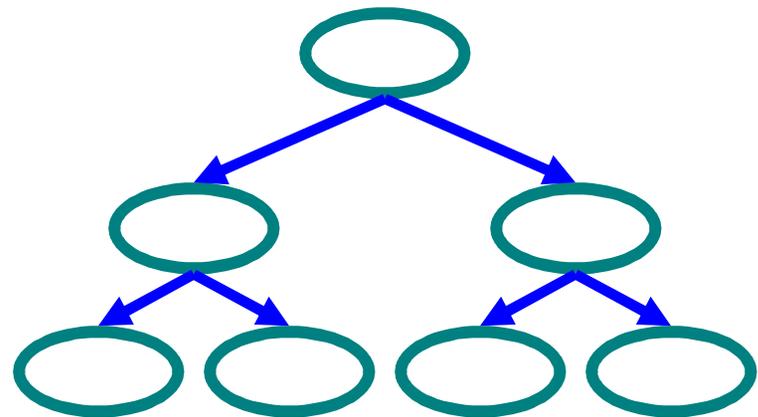
## Tree Data Structures

# Trees Data Structures

- Tree
  - Nodes
  - Each node can have 0 or more **children**
  - A node can have at most one **parent**
- Binary tree
  - Tree with 0–2 children per node



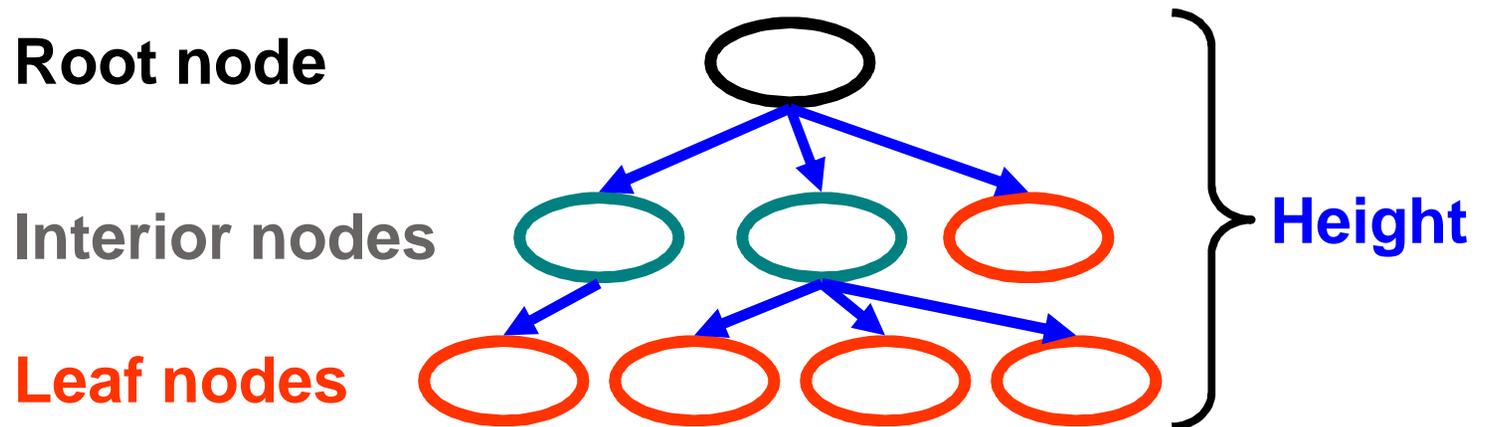
**Tree**



**Binary Tree**

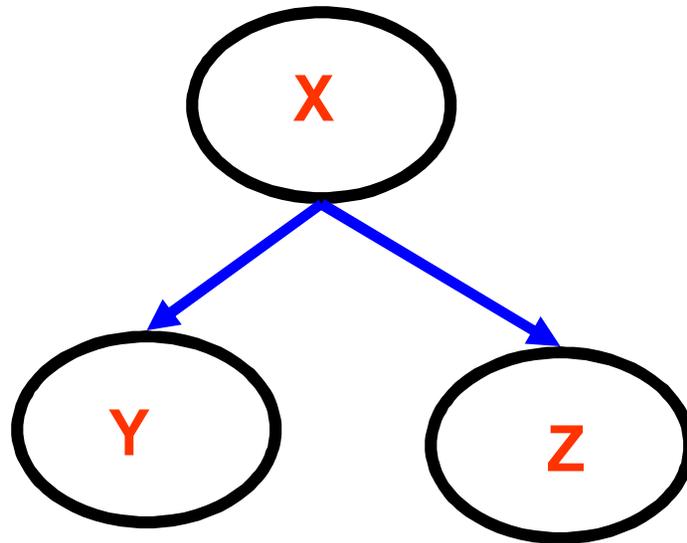
# Trees

- Terminology
  - Root  $\Rightarrow$  no parent
  - Leaf  $\Rightarrow$  no child
  - Interior  $\Rightarrow$  non-leaf
  - Height  $\Rightarrow$  distance from root to leaf



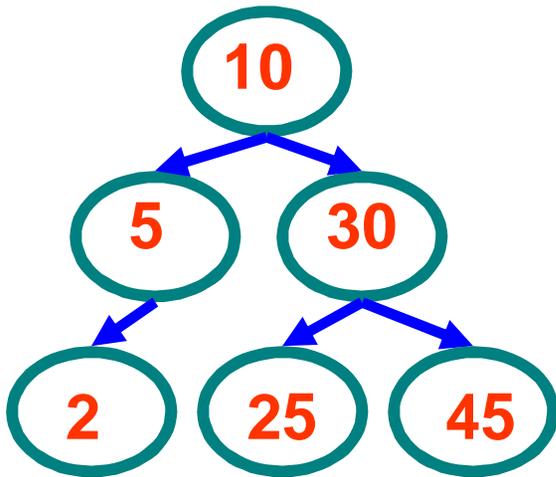
# Binary Search Trees

- Key property
  - Value at node
    - Smaller values in left subtree
    - Larger values in right subtree
  - Example
    - $X > Y$
    - $X < Z$

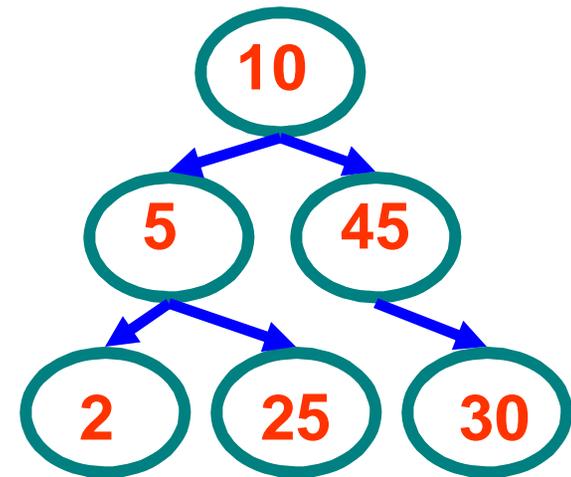
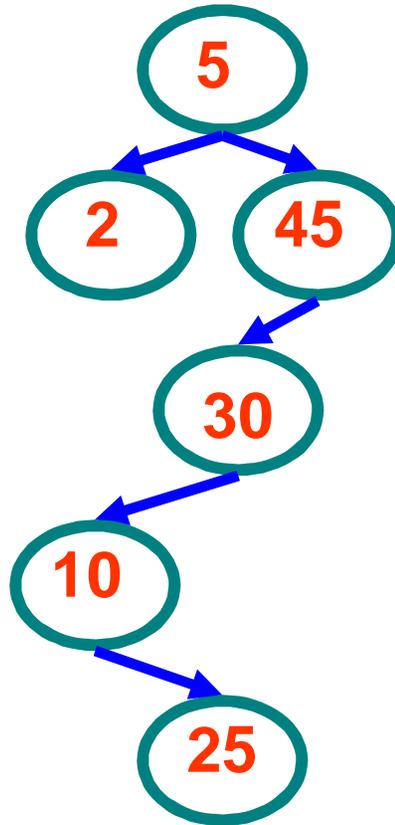


# Binary Search Trees

- Examples



**Binary  
search trees**



**Not a binary  
search tree**

# Binary Tree Implementation

```
Class Node {  
    int data; // Could be int, a class, etc  
    Node *left, *right; // null if empty  
  
    void insert ( int data ) { ... }  
    void delete ( int data ) { ... }  
    Node *find ( int data ) { ... }  
    ...  
}
```

# Iterative Search of Binary Tree

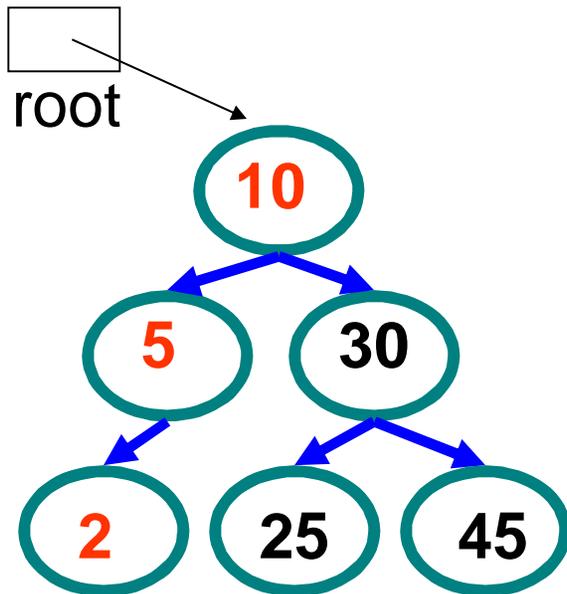
```
Node *Find( Node *n, int key) {  
    while (n != NULL) {  
        if (n->data == key)           // Found it  
            return n;  
        if (n->data > key)             // In left subtree  
            n = n->left;  
        else                           // In right subtree  
            n = n->right;  
    }  
    return null;  
}  
Node * n = Find( root, 5);
```

# Recursive Search of Binary Tree

```
Node *Find( Node *n, int key) {  
    if (n == NULL)           // Not found  
        return( n );  
    else if (n->data == key)  // Found it  
        return( n );  
    else if (n->data > key)   // In left subtree  
        return Find( n->left, key );  
    else                       // In right subtree  
        return Find( n->right, key );  
}  
Node * n = Find( root, 5);
```

# Example Binary Searches

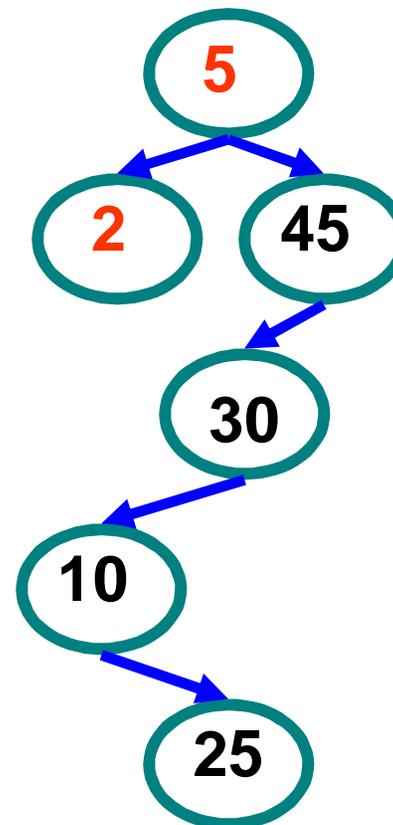
- Find ( root, 2 )



$10 > 2$ , left

$5 > 2$ , left

$2 = 2$ , found

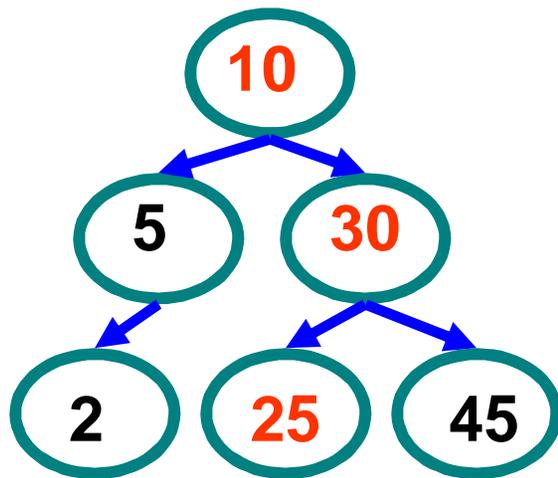


$5 > 2$ , left

$2 = 2$ , found

# Example Binary Searches

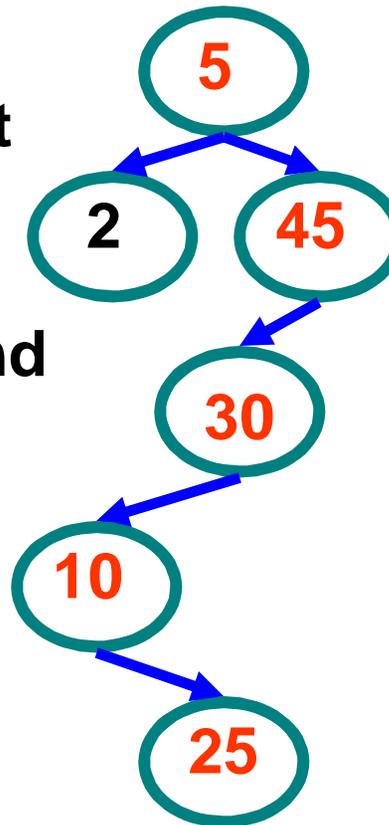
- Find (root, 25 )



$10 < 25$ , right

$30 > 25$ , left

$25 = 25$ , found



$5 < 25$ , right

$45 > 25$ , left

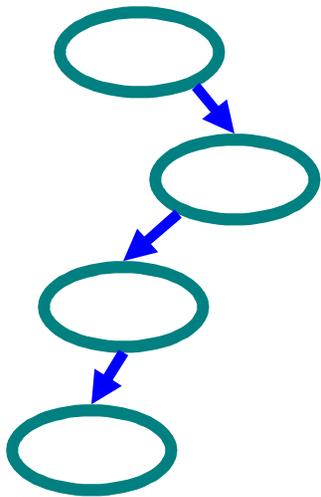
$30 > 25$ , left

$10 < 25$ , right

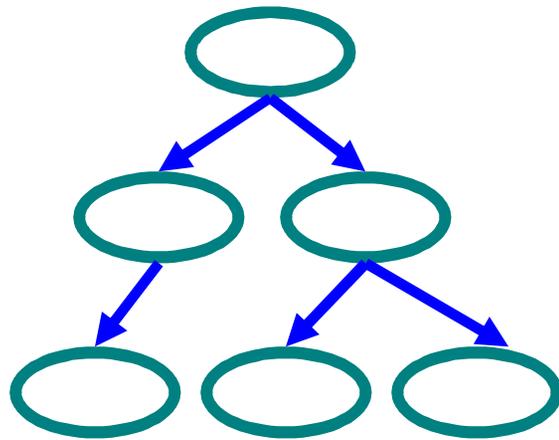
$25 = 25$ , found

# Types of Binary Trees

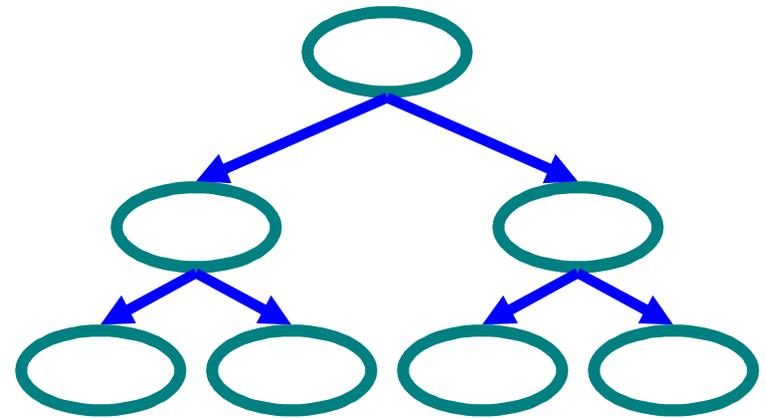
- Degenerate – only one child
- Complete – always two children
- Balanced – “mostly” two children
  - more formal definitions exist, above are intuitive ideas



**Degenerate  
binary tree**



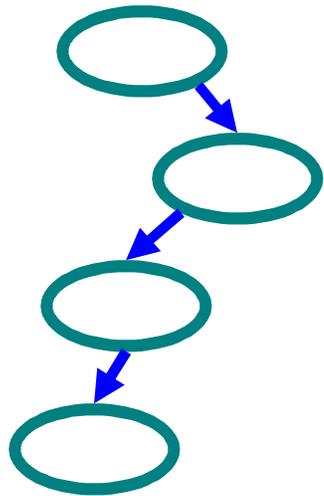
**Balanced  
binary tree**



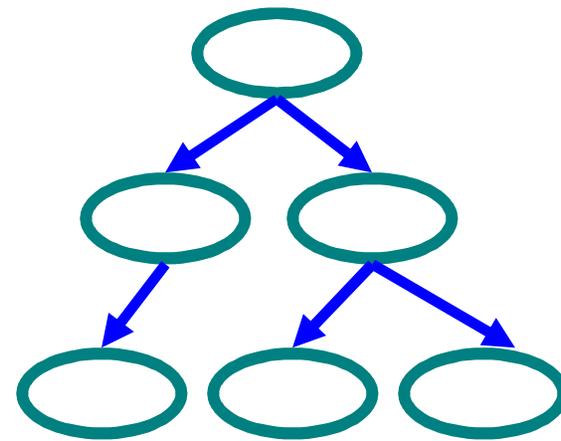
**Complete  
binary tree**

# Binary Trees Properties

- Degenerate
  - Height =  $O(n)$  for  $n$  nodes
  - Similar to linked list
- Balanced
  - Height =  $O(\log(n))$  for  $n$  nodes
  - Useful for searches



**Degenerate  
binary tree**



**Balanced  
binary tree**

# Binary Search Properties

- Time of search
  - Proportional to height of tree
  - Balanced binary tree
    - $O(\log(n))$  time
  - Degenerate tree
    - $O(n)$  time
    - Like searching linked list / unsorted array

# Binary Search Tree Construction

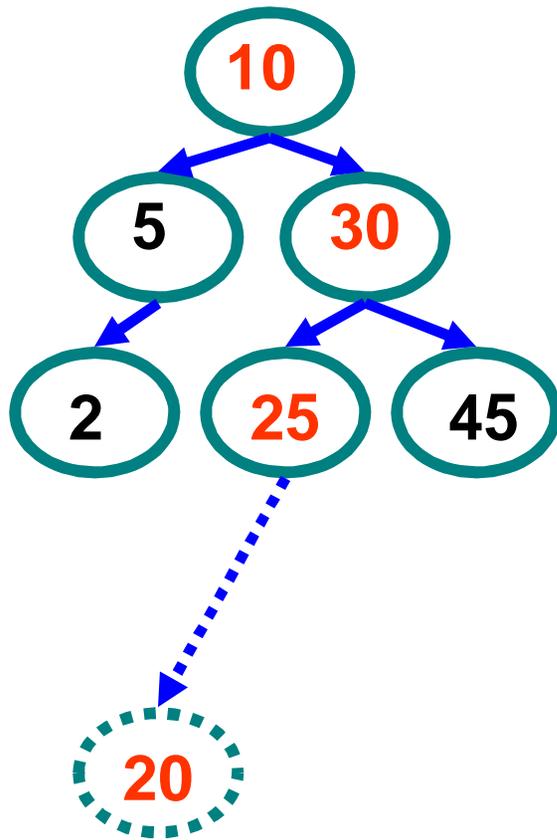
- How to build & maintain binary trees?
  - Insertion
  - Deletion
- Maintain key property (invariant)
  - Smaller values in left subtree
  - Larger values in right subtree

# Binary Search Tree – Insertion

- Algorithm
  1. Perform search for value X
  2. Search will end at node Y (if X not in tree)
  3. If  $X < Y$ , insert new leaf X as new left subtree for Y
  4. If  $X > Y$ , insert new leaf X as new right subtree for Y
- Observations
  - $O(\log(n))$  operation for balanced tree
  - Insertions may unbalance tree

# Example Insertion

- Insert ( 20 )



**10 < 20, right**

**30 > 20, left**

**25 > 20, left**

**Insert 20 on left**

# Binary Search Tree – Deletion

- Algorithm

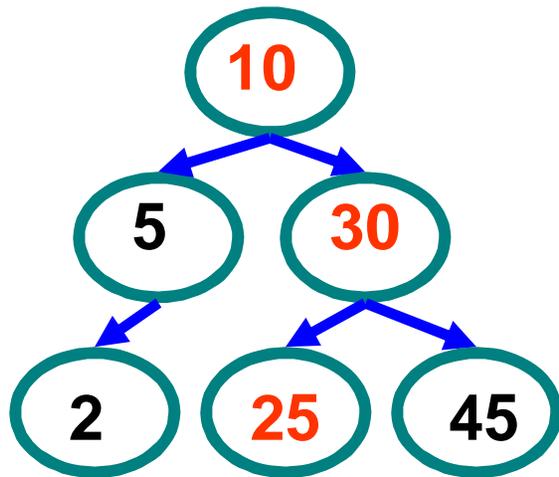
1. Perform search for value X
2. If X is a leaf, delete X
3. Else // must delete internal node
  - a) Replace with largest value Y on left subtree  
OR smallest value Z on right subtree
  - b) **Delete** replacement value (Y or Z) from subtree

- Observation

- $O(\log(n))$  operation for balanced tree
- Deletions may unbalance tree

# Example Deletion (Leaf)

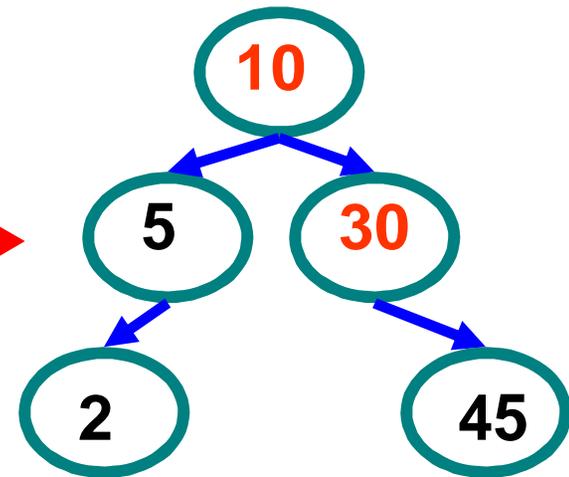
- Delete ( 25 )



$10 < 25$ , right

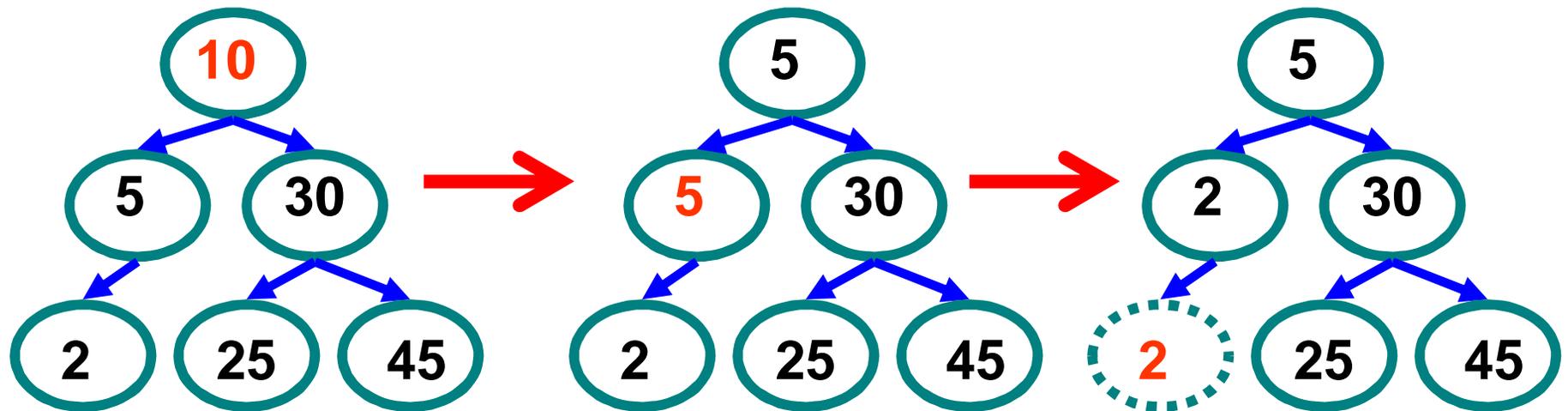
$30 > 25$ , left

$25 = 25$ , delete



# Example Deletion (Internal Node)

- Delete ( 10 )



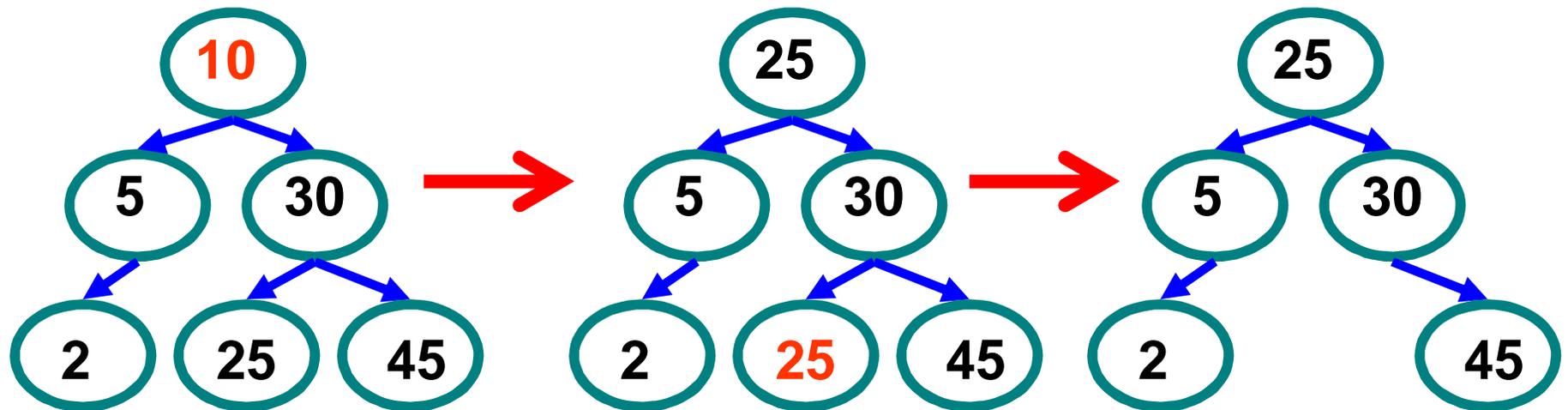
Replacing 10  
with **largest**  
value in left  
subtree

Replacing 5  
with **largest**  
value in left  
subtree

Deleting leaf

# Example Deletion (Internal Node)

- Delete ( 10 )



Replacing 10  
with **smallest**  
value in right  
subtree

Deleting leaf

Resulting tree

# Balanced Search Trees

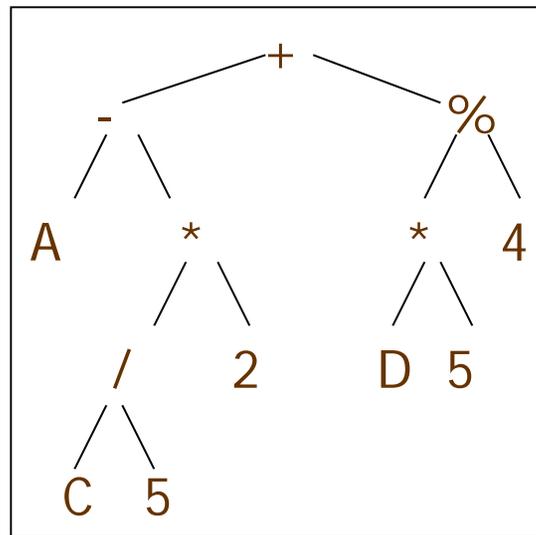
- Kinds of balanced binary search trees
  - height balanced vs. weight balanced
  - “Tree rotations” used to maintain balance on insert/delete
- Non-binary search trees
  - 2/3 trees
    - each internal node has 2 or 3 children
    - all leaves at same depth (height balanced)
  - B-trees
    - Generalization of 2/3 trees
    - Each internal node has between  $k/2$  and  $k$  children
      - Each node has an array of pointers to children
    - Widely used in databases

# Other (Non-Search) Trees

- Parse trees
  - Convert from textual representation to tree representation
  - Textual program to tree
    - Used extensively in compilers
  - Tree representation of data
    - E.g. HTML data can be represented as a tree
      - called DOM (Document Object Model) tree
    - XML
      - Like HTML, but used to represent data
      - Tree structured

# Parse Trees

- Expressions, programs, etc can be represented by tree structures
  - E.g. Arithmetic Expression Tree
  - $A - (C / 5 * 2) + (D * 5 \% 4)$



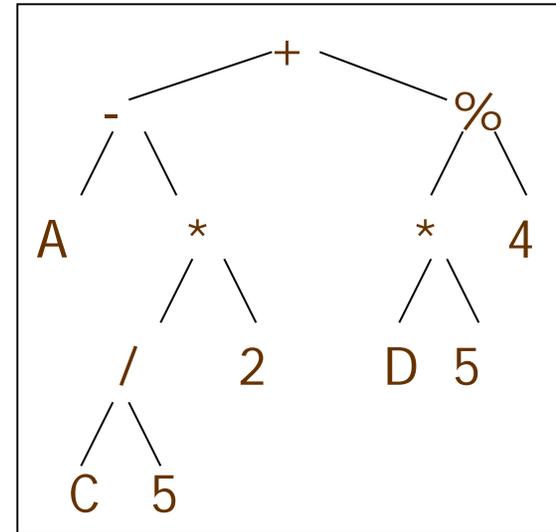
# Tree Traversal

- Goal: visit every node of a tree
- **in-order** traversal

```
void Node::inOrder () {  
    if (left != NULL) {  
        cout << "("; left->inOrder(); cout << ")";  
    }  
    cout << data << endl;  
    if (right != NULL) right->inOrder()  
}
```

Output: **A - C / 5 \* 2 + D \* 5 % 4**

To disambiguate: print brackets



# Tree Traversal (contd.)

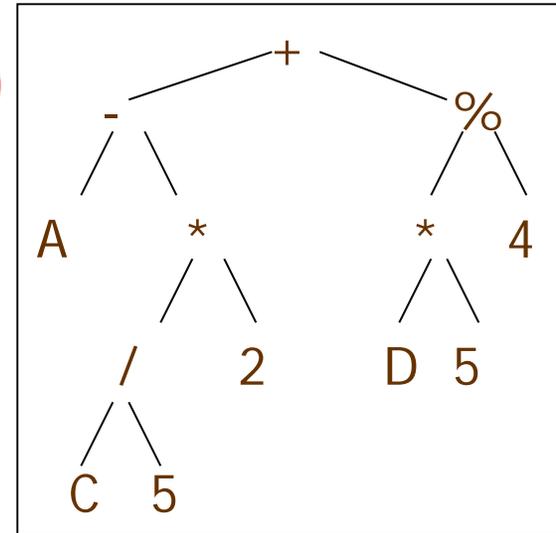
- **pre-order** and **post-order**:

```
void Node::preOrder () {  
    cout << data << endl;  
    if (left != NULL) left->preOrder ();  
    if (right != NULL) right->preOrder ();  
}
```

Output: + - A \* / C 5 2 % \* D 5 4

```
void Node::postOrder () {  
    if (left != NULL) left->preOrder ();  
    if (right != NULL) right->preOrder ();  
    cout << data << endl;  
}
```

Output: A C 5 / 2 \* - D 5 \* 4 % +



# XML

- Data Representation

- E.g.

- <dependency>

- <object>sample1.o</object>

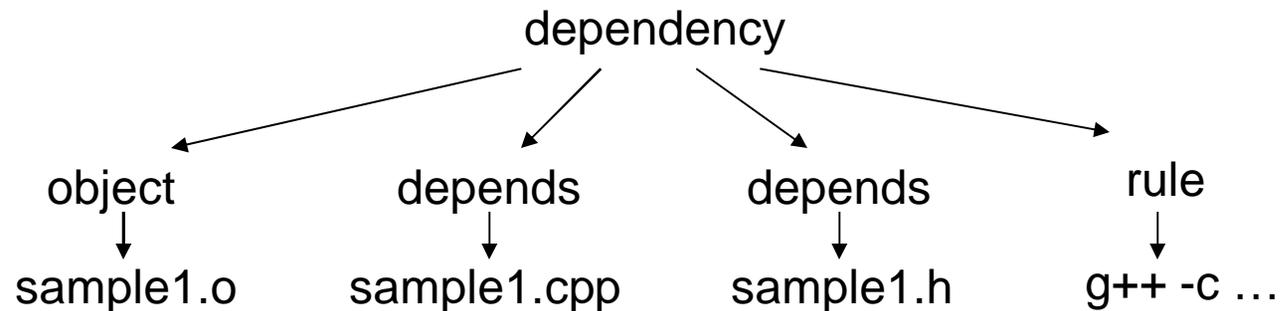
- <depends>sample1.cpp</depends>

- <depends>sample1.h</depends>

- <rule>g++ -c sample1.cpp</rule>

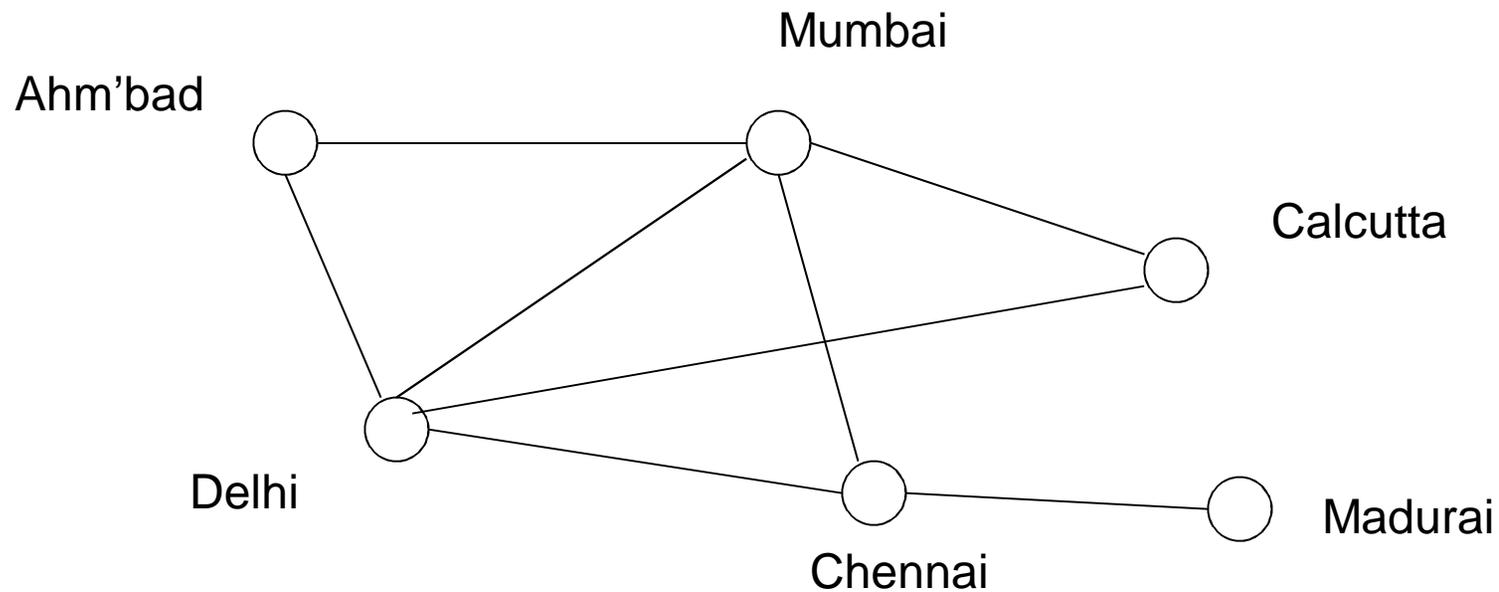
- </dependency>

- Tree representation



# Graph Data Structures

- E.g: Airline networks, road networks, electrical circuits
- Nodes and Edges
- E.g. representation: class Node
  - Stores name
  - stores pointers to all adjacent nodes
    - i,e. edge == pointer
    - To store multiple pointers: use array or linked list



End of Chapter